

# Topic 3:

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## Quantification

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## Propositions With Variables (1 / 2)

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Propositions are static; variables are not allowed. But ...

**Definition: Predicate (a.k.a. Propositional Function)**

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**Example(s):**

# Propositions With Variables (2 / 2)

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## Definition: Domain (a.k.a. Universe) of Discourse

## Example(s):

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## Quantification

Idea: Establish truth of predicates over sets of values.

Two common generalizations:

**Note:** Do not use the book's unusual  $\exists!x$  notation.

# Evaluating Quantified Predicates (1 / 2)

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## 1. Universally Quantified Predicates

**Example(s):**



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# Evaluating Quantified Predicates (2 / 2)

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## 2. Existentially Quantified Predicates

**Example(s):**



## Evaluating Mixed Quantifications (1 / 2)

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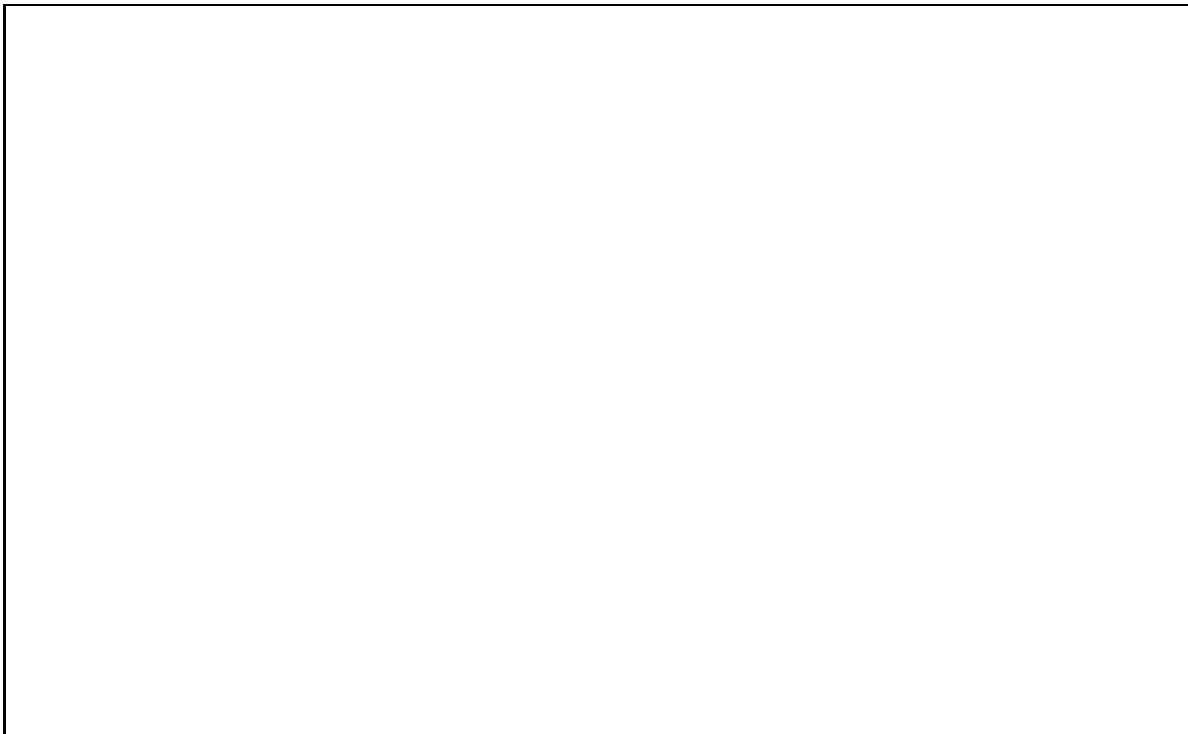
First: Distinguishing  $\exists x \forall y S(x, y)$  from  $\forall i \exists k T(i, k)$  :

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## Evaluating Mixed Quantifications (2 / 2)

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**Example(s):**



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## Example: Universal Quantification (1 / 5)

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Consider this conversational English statement:

All of McCann's students are geniuses.

How can we express that statement in logic notation?

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## Example: Universal Quantification (2 / 5)

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Attempt #2: All of McCann's students are geniuses.  $\rightarrow$  Logic

## Example: Universal Quantification (3 / 5)

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Attempt #3: All of McCann's students are geniuses.  $\rightarrow$  Logic

Let  $P(x)$  : Student  $x$  is a genius,  $x \in \text{People}$

## Example: Universal Quantification (4 / 5)

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Attempt #4: All of McCann's students are geniuses.  $\rightarrow$  Logic

Let  $P(x)$  : Student  $x$  is a genius,  $x \in \text{People}$

Let  $M(x)$  :  $x$  is enrolled in one of McCann's classes,  $x \in \text{People}$

## Example: Universal Quantification (5 / 5)

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Attempt #5: All of McCann's students are geniuses.  $\rightarrow$  Logic

Let  $P(x)$  : Student  $x$  is a genius,  $x \in \text{People}$

Let  $M(x)$  :  $x$  is enrolled in one of McCann's classes,  $x \in \text{People}$

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## Implicit Quantification

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The “all” can be implicit in the English statement.

**Example(s):**

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## Example: Existential Quantification

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Consider this conversational English statement:

At least one towel is dirty.

How can we express that statement in logic notation?

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## Another Example: Existential Quantification

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Express this more specific statement in logic:

Some of the blue guest towels are dirty.

Let  $D(x)$  :  $x$  is dirty,  $x \in \text{Towels}$



# Yet Another Example: Quantification

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Now express this statement in logic:

Every last one of the blue guest towels are dirty!

Let  $B(x)$  :  $x$  is blue,  $x \in \text{Towels}$

Let  $G(x)$  :  $x$  is used by guests,  $x \in \text{Towels}$

Let  $D(x)$  :  $x$  is dirty,  $x \in \text{Towels}$

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## Free vs. Bound Variables

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**Definition: Bound Variable**

**Definition: Free (a.k.a. Unbound) Variable**

Other examples of 'binding' in CS:

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# Negations of Quantified Expressions

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Remember De Morgan's Laws for propositions? Well, ...

## Definition: Generalized De Morgan's Laws

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## Demonstration: $\overline{\forall x P(x)} \equiv \exists x \overline{P(x)}$ (1 / 2)

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Let  $S(x) : x < 4, x \in \mathbb{Z}$

The expression  $\forall x S(x), x \in \{1, 2, 3\}$  is true.

Equivalently,  $\overline{\forall x S(x)}$  is false.

$$\forall x S(x) \equiv S(1) \wedge S(2) \wedge S(3) \quad \text{so ...}$$

$$\begin{aligned} \overline{\forall x S(x)} &\equiv \overline{S(1) \wedge S(2) \wedge S(3)} \\ &\equiv \overline{S(1)} \vee \overline{S(2)} \vee \overline{S(3)} \quad \text{[De Morgan, 2x]} \end{aligned}$$

(Remember:  $\overline{S(1)} \vee \overline{S(2)} \vee \overline{S(3)}$  is still false.)

(Continues ...)

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## Demonstration: $\overline{\forall x P(x)} \equiv \exists x \overline{P(x)}$ (2 / 2)

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For  $\overline{S(1)} \vee \overline{S(2)} \vee \overline{S(3)}$  to be false, each term must be false; that is, no  $\overline{S(x)}$  is true (or all  $\overline{S(x)}$  are false).

It follows that the expression  $\exists x \overline{S(x)}$  must be false, completing the demonstration.

Example(s):

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## Expressing “Exactly one . . .” Statements (1 / 3)

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Consider this conversational (& correct!) English statement:

Only one citizen of North Dakota is a member of  
the U.S. House of Representatives.

And consider this awkward but useful rewording:

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## Expressing “Exactly one . . .” Statements (2 / 3)

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That rewording is useful because it can be directly expressed logically:

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## Expressing “Exactly one . . .” Statements (3 / 3)

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A lingering problem:

The domain (“Citizens of North Dakota”) is too specific.

Solution: Add a predicate . . .but what, and where?

## Expressing “Exactly two ...” Statements (1 / 3)

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Key observation:

**Question:** Can you say this in ‘awkward English’?

**Exactly two citizens of North Dakota are U.S. Senators.**

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## Expressing “Exactly two ...” Statements (2 / 3)

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Consider the two halves separately. Given:

$S(x)$  :  $x$  is a U.S. Senator,  $x \in \text{People}$

1. “At least two citizens of North Dakota are U.S. Senators”
  
  
  
  
  
  
  
  
  
  
2. “At most two citizens of North Dakota are U.S. Senators”

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## Expressing “Exactly two ...” Statements (3 / 3)

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Finally, AND together

$$\exists x \exists y (S(x) \wedge S(y) \wedge (x \neq y))$$

and

$$\forall x \forall y \forall z ((S(x) \wedge S(y) \wedge S(z)) \\ \rightarrow (x = y \vee y = z \vee x = z)):$$