

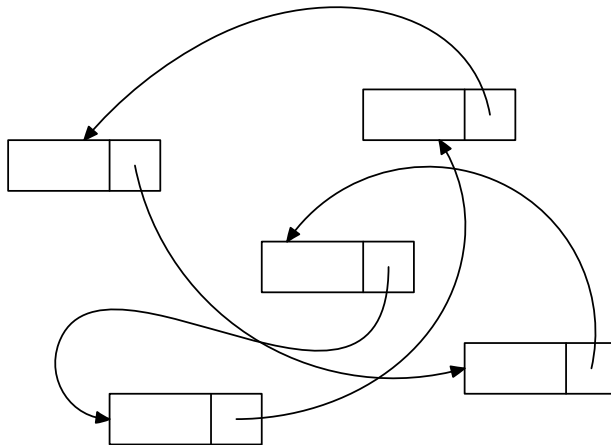
Topic 1:

Review Material

(or: Let's Spackle Some Knowledge Holes!)

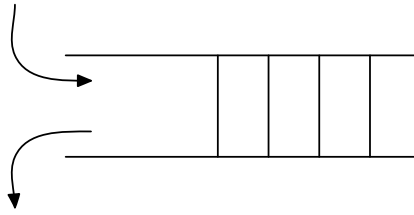
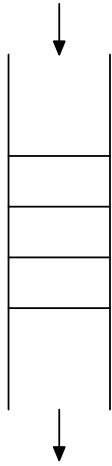
Review – CSc 345 v1.0 (McCann) – p. 1/33

What is This, and What is Missing?



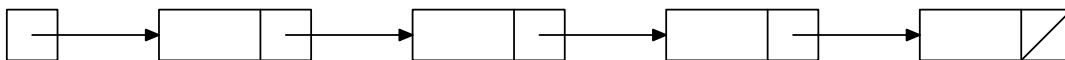
Review – CSc 345 v1.0 (McCann) – p. 2/33

What are these Data Structures?



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Singly-Linked List (SLL) Review (1 / 2)



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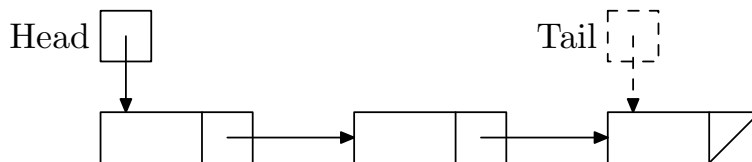
Singly-Linked List (SLL) Review (2 / 2)

What are some common SLL operations?

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SLLs: Tail Reference or no Tail Reference?

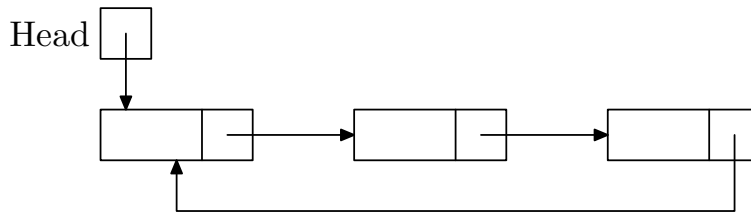
When would a Tail reference be helpful? Or annoying?



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Circularly-Linked Lists (CLLs)

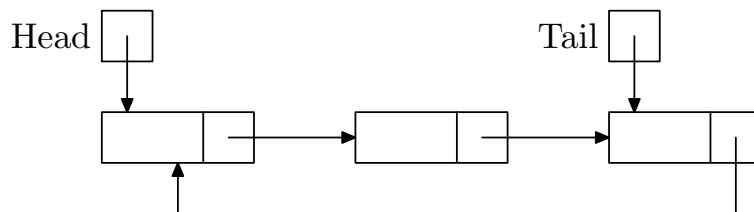
Idea: Allow easy travel from last node to first.



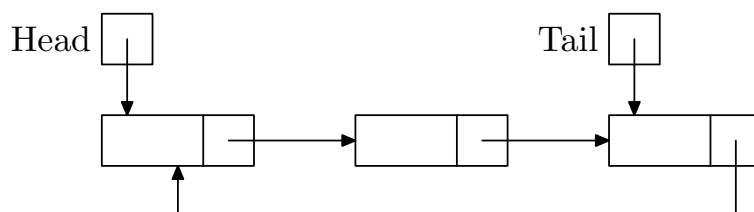
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CLLs: Are Tail References Helpful?

Append



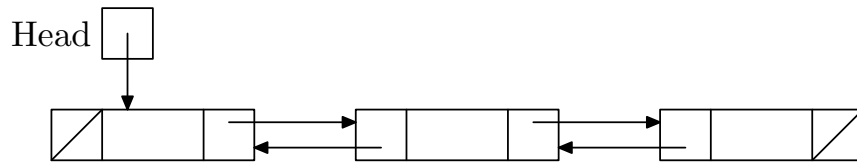
Prepend



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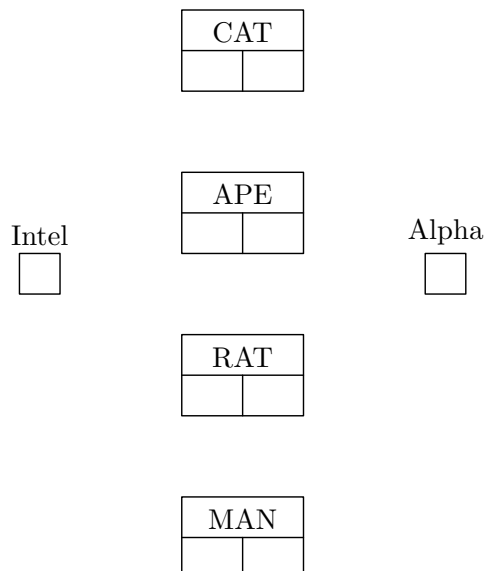
Doubly-Linked Lists (DLLs)

An extra reference per node allows two-way travel.

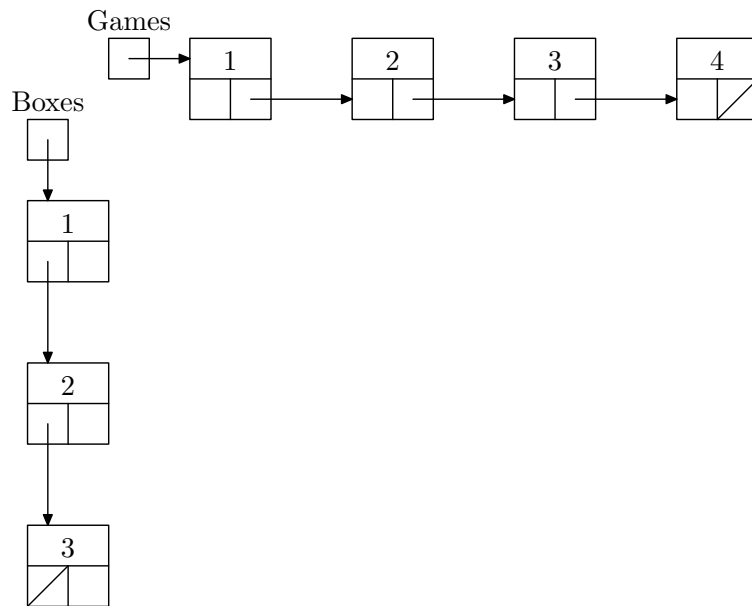


MultiLists

Here, each node is part of two distinct lists of the same data.



Orthogonal Lists



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1D Array Storage

Array elements are stored *contiguously* in memory.

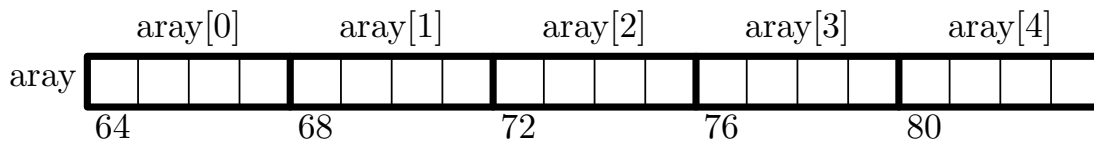
Example(s):



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Where in Memory is `array[index]`? (1 / 3)

We need to know a few things first:

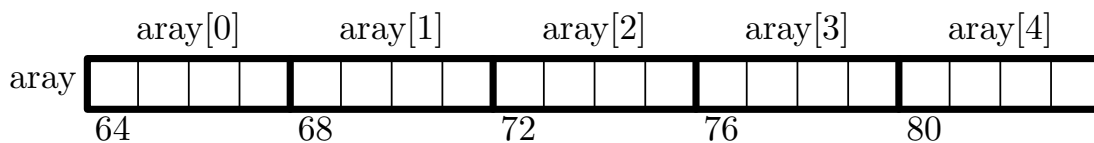


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-
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Where in Memory is `array[index]`? (2 / 3)

To reach index i 's element, we have to 'jump' over i elements.



Recall: `esize = 4`, `base = 64`, `index = 3`

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Where in Memory is `array[index]`? (3 / 3)

OK, great,

```
address = base + index * esize
```

but ... so what?

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2D Array Allocation in Java

There are more options than you might guess. Two of them:

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2D Array Storage (1 / 5)

How can we store a 2D structure in 1D memory? Two options:

	0	1	2	3
0	6	0	8	2
1	1	5	3	6
2	9	4	7	1

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2D Array Storage (2 / 5)

Consider locating 4 in our 2D array:

	0	1	2	3
0				
1				
2		4		

RMO:

							4	
--	--	--	--	--	--	--	---	--

CMO:

				4				
--	--	--	--	---	--	--	--	--

Our 2D array address calculation is in two parts:

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2D Array Storage (3 / 5)

For our 1D calculation, we needed:

- base address
- esize
- index

What information do we need for the 2D calculation?

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2D Array Storage (4 / 5)

Equation time! (I'll do RMO; you can do CMO on your own.)

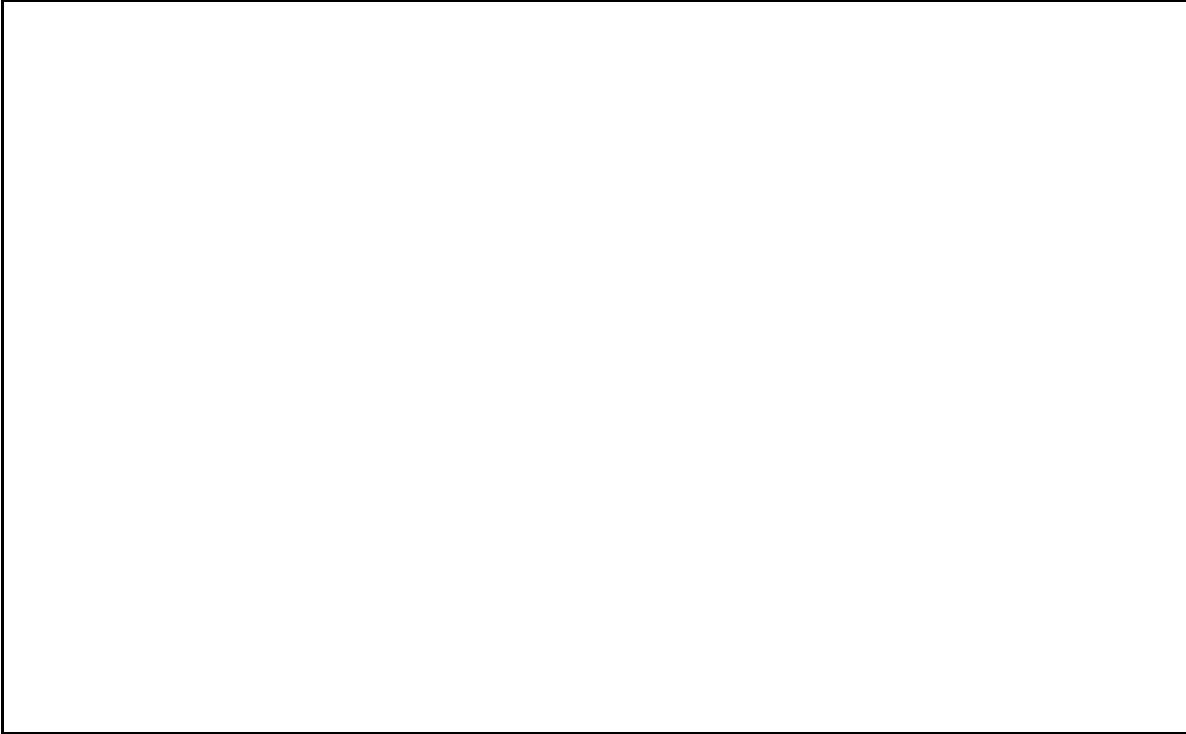
twod

									4		
--	--	--	--	--	--	--	--	--	---	--	--

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2D Array Storage (5 / 5)

Example(s):



Review – CSc 345 v1.0 (McCann) – p. 21/33

What About 3D Arrays?

The same basic idea, just one more dimension!

Think about slicing pre-cROUTONS
from a loaf of bread. \implies



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*n*D Arrays in OO Languages (1 / 3)

In object-oriented (OO) languages, arrays are objects.

A 1D array object is a contiguous collection of references to data objects.

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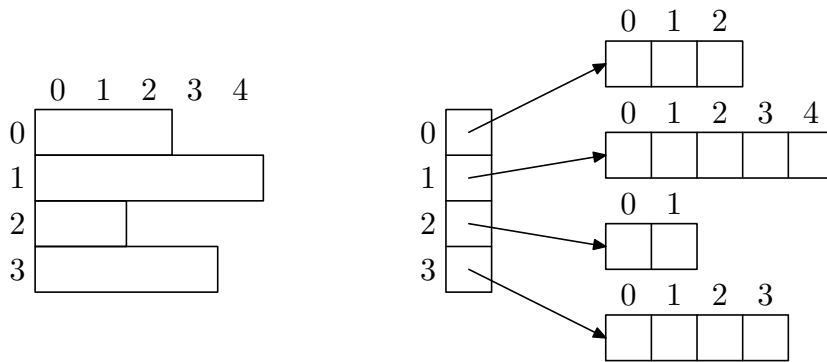
*n*D Arrays in OO Languages (2 / 3)

Because each row is a distinct object, they can ...

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*n*D Arrays in OO Languages (3 / 3)

How do we declare & allocate such arrays in Java?



Recursion Review

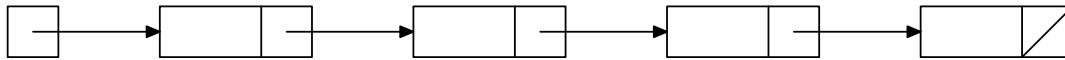
Want to solve a problem recursively? Try to answer both of these questions:

1.

2.

Simple Recursion Example #1 (1 / 2)

Task: Print the content of an SLL front to rear.



Q1: What's somewhat simpler than printing an SLL of n items front to rear?

Q2: How does that help print the list of n items front to rear?

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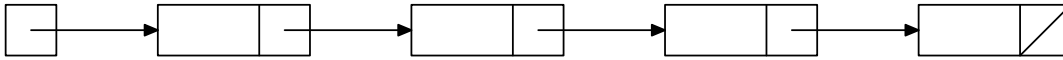
Simple Recursion Example #1 (2 / 2)

Let's turn our answers into a pseudocode algorithm:

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Simple Recursion Example #2 (1 / 2)

Task: Print the content of an SLL rear to front.



Q1: What's somewhat simpler than printing an SLL of n items rear to front?

Q2: How does that help print the list of n items rear to front?

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Simple Recursion Example #2 (2 / 2)

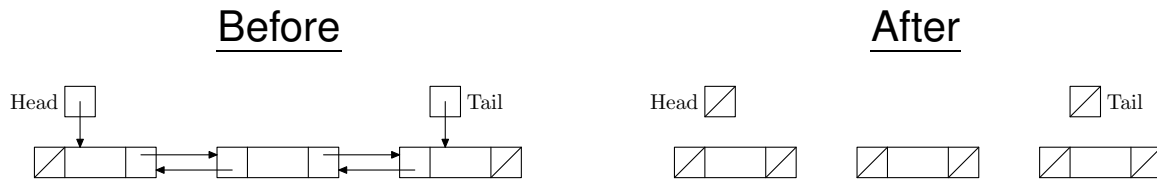
There's just one change from the 'front to rear' version:

```
1 subprogram printLL (given: head) returns nothing
2   if head is null:
3     return
4   else:
5     call printLL with head's successor
6     print head's data
7   end if
8 end subprogram
```

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Simple Recursion Example #3 (1 / 3)

Task: Totally unlink a DLL.



Q1: What's somewhat simpler than unlinking a DLL of n items?

Q2: How does that help unlink the DLL of n items?

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Simple Recursion Example #3 (2 / 3)

In pseudocode:

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Simple Recursion Example #3 (3 / 3)

We need to set Head and Tail to null . . . but when and where?

Pseudocode: