

Topic 4:

Syntax

Q: How do you comfort a grammar expert?

A: “There, their, they’re”

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Syntax v. Semantics

Definition: Syntax

.....

Semantics are related but distinct:

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Tokens and Lexemes

Definition: Token

.....

Definition: Lexeme

Sometimes a token and a lexeme are the same:

Token:

Lexeme:

But usually not:

Token:

Lexemes:

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Lexical Structure, Scanning, and Parsing

Definition: Lexical Structure

.....

During language translation:

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Identifying Tokens (1 / 2)

Where does one lexeme stop and the next start?

Definition: Principle of Longest Substring

Example from Java:

```
int i=0;  doi++;  while (i<10);
```

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Identifying Tokens (2 / 2)

Example from FORTRAN 66:

```
DO 99 I=1,10    v.    DO 99 I=1.10
```

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Context–Free Grammars & Backus–Naur Form (1 / 2)

In the mid–1950s, linguist **Noam Chomsky** defined a relationship, known as the **Chomsky Hierarchy**, between four classes of languages:

Type 0:

Type 1:

Type 2:

Type 3:

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Context–Free Grammars & Backus–Naur Form (2 / 2)

Definition: Grammar

.....

John Backus and **Peter Naur** defined and refined a notation for describing programming language syntax while developing the language **ALGOL**.

The connection:

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Regular and Context–Free Languages (1 / 2)

Regular languages can be expressed by regular expressions.

Be Aware! — Regular expressions, as provided in programming languages, add additional features that go beyond true regular languages.

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Regular and Context–Free Languages (2 / 2)

Context–free languages are defined by context–free grammars (CFGs).

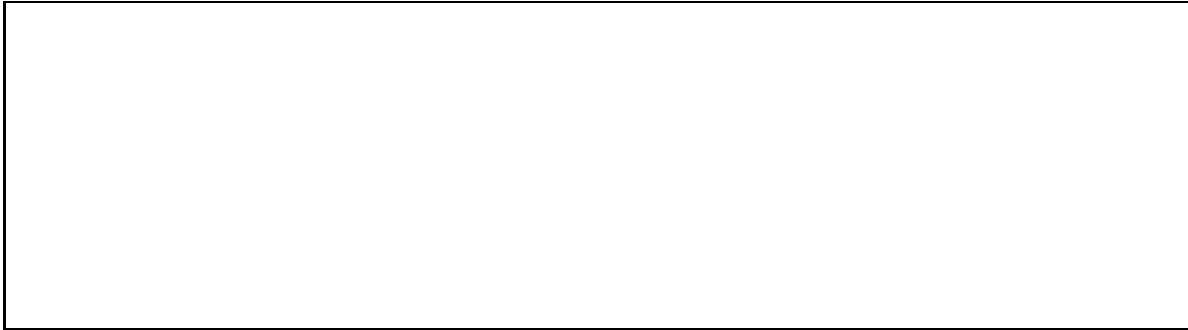
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BNF Productions (1 / 3)

Note: This style of writing BNF productions (rules) differs from our textbook's style. Why not be consistent? Italics are hard to write by hand!

Productions have this form:

Example(s): BNF for a C language assignment statement

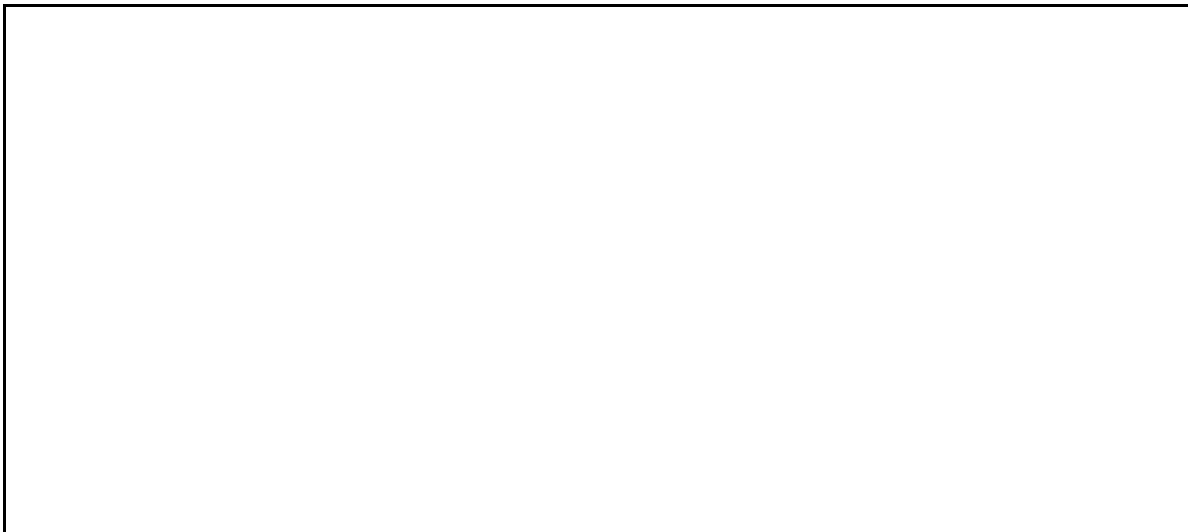


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BNF Productions (2 / 3)

Non-terminals may have multiple definitions.

Example(s): Java IF statements



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BNF Productions (3 / 3)

BNF rules look simple, but hide a lot of power!

Example(s): BNF can describe comma-separated lists



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A Larger BNF Example

Consider this grammar, in which `<block>` is the start symbol.

An expression of the language defined by this grammar:

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Constructing a Derivation (1 / 2)

Definition: Derivation

.....
.....

One possible derivation algorithm ('RHS' $\equiv$ Right-Hand Side):

Write down the start symbol's rule

Repeat until all non-terminals have been removed:

Find the left-most non-terminal on the RHS of the rule

Replace it with the RHS of one of its rules

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Constructing a Derivation (2 / 2)

Note: The symbol $\Rightarrow$ means “derives.”

Let's derive our expression: **start I = J * K stop**

<block> $\Rightarrow$ start <stmt-list> stop

$\Rightarrow$

$\Rightarrow$

$\Rightarrow$

$\Rightarrow$

$\Rightarrow$

$\Rightarrow$

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Parse Trees (a.k.a. Concrete Syntax Trees)

Here's a hierarchical view of the derivation process:

(Recall: **start I = J * K stop**)

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Abstract Syntax Trees

Most of the non-terminals in Concrete Parse Trees aren't very helpful, so why keep drawing them?

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Ambiguous Grammars (1 / 6)

Consider this enhanced version of our assignment statement grammar:

$$\langle \text{stmt} \rangle \rightarrow \langle \text{var} \rangle = \langle \text{expr} \rangle$$
$$\langle \text{var} \rangle \rightarrow \text{I} \mid \text{J} \mid \text{K}$$

We'll use it to derive this expression:

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Ambiguous Grammars (2 / 6)

Problem! That grammar can derive two syntax trees:

$$\text{K} = \text{I} * \text{J} / \text{K}$$

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Ambiguous Grammars (3 / 6)

How can we remove this ambiguity?

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Ambiguous Grammars (4 / 6)

Let's look at a rewriting option for this grammar:

`<stmt> → <var> = <expr>`

`<var> → I | J | K`

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Ambiguous Grammars (5 / 6)

$I / J * K$

v.

$I / (J * K)$

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Ambiguous Grammars (6 / 6)

Notes:

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Extended BNF (EBNF) (1 / 4)

Ref: Wirth, Niklaus. "What Can We Do about the Unnecessary Diversity of Notation for Syntactic Definitions?" CACM, 11/1977.

We've seen BNF notations like this:

$\langle \text{expr} \rangle \rightarrow \langle \text{expr} \rangle * \langle \text{term} \rangle \mid \langle \text{term} \rangle$

Question:

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Extended BNF (EBNF) (2 / 4)

Added notations (but not power!) in EBNF include:

$\{ \}$ —

$[]$ —

$()$ —

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Extended BNF (EBNF) (3 / 4)

Example(s):

Wirth's language **Pascal** optionally lists files on the first line:

```
program foo (input, output) ;
```

We can express this in BNF:

Or more compactly in EBNF:

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Extended BNF (EBNF) (4 / 4)

Problem! EBNF parse trees are sometimes ambiguous.

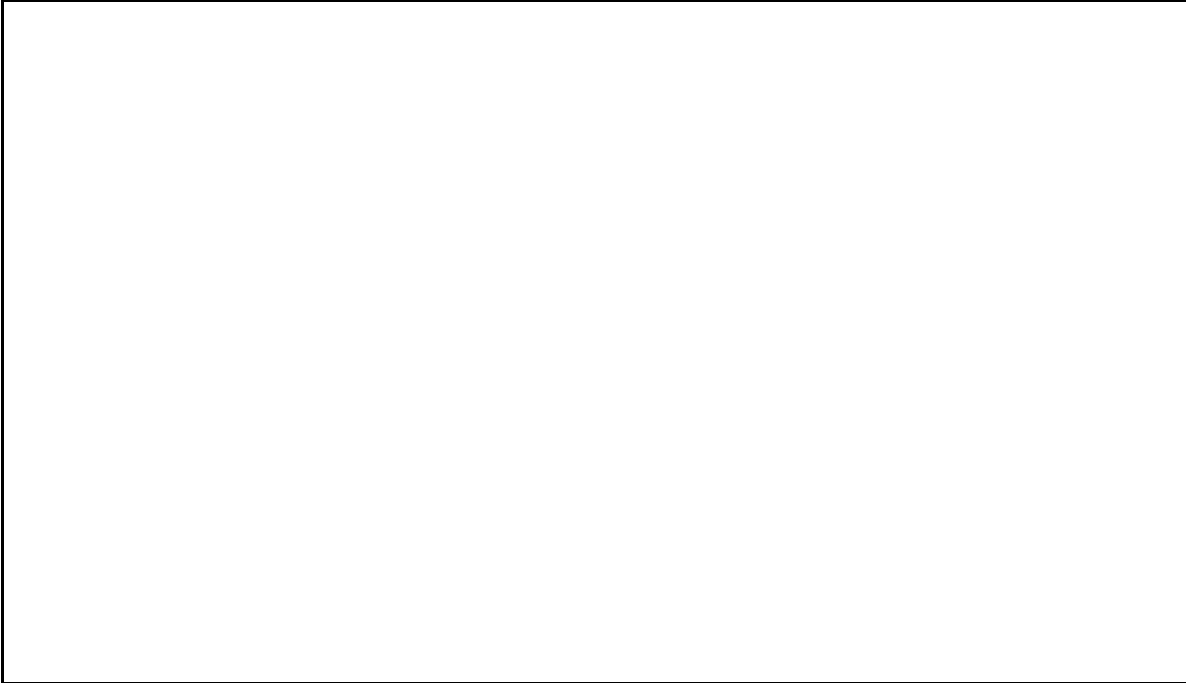
Example(s):

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Syntax Diagrams (a.k.a. Railroad Diagrams)

Looking for a visualization of (E)BNF?

Example(s):

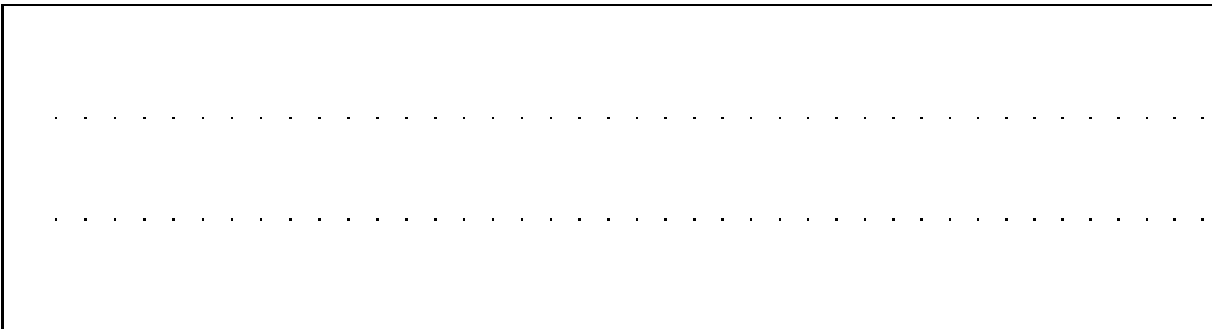


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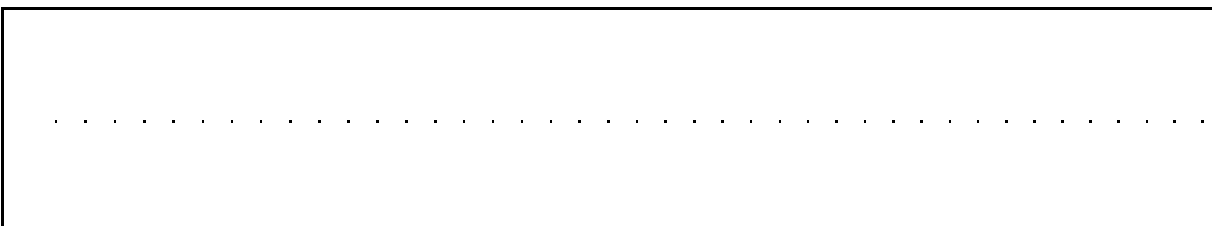
Parsing (1 / 3)

Which programs test code against a language's expectations?

Definition: Parser



Definition: Recognizer



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Parsing (2 / 3)

There are two major types of parsers.

- 1.

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Parsing (3 / 3)

- 2.

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Concluding Notes on Syntax

Which program checks that the tokens are correctly constructed?

CFGs are capable of more than mere syntax-checking