Extending OpenMP to Facilitate Loop Optimization

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Outline

- Limitations of OpenMP.
- Motivation for sequential transformations in OpenMP.
- Proposed Optimizations.
 - o Unroll
 - Fuse
 - Tile
- Composition of loop transformations.
- Interaction with other OpenMP constructs.
- Discussion.

Limitations

OpenMP Does:

- Loop Parallelism
- Task Parallelism
- Accelerator Offloading

OpenMP Doesn't:

- Basic loop transformations
 - \circ Unrolling
 - Fusion
 - \circ Tiling

Why should OpenMP provide sequential loop optimizations?

- 1. Provide standard for portable sequential loop optimization directives and semantics.
- 2. Sequential optimizations improve parallel code performance.
- 3. OpenMP semantics obscure/prevent them from being automatically performed by the compiler.

Performance Improvements with Loop Unroll

CLOMPK (CORAL2 benchmark)



Possible Unroll Syntax and Semantics Example

```
#pragma omp for unroll( 2 ) schedule( static, 1 ) nowait
for( int i = 0; i < n; i += 1 )</pre>
 A[i] += B[i] * c;
#pragma omp for schedule( static, 1 ) nowait
for( int i = 0; i < n; i += 2){
 A[i] += B[i] * c;
 A[i+1] += B[i+1] * c;
}
```

Context:

• Data sharing between loop nests are an opportunity to exploit caching and reduce temporary storage.

Problem:

- Manual modifications difficult.
- Automatic transformation requires complex analysis.

Solution:

- Provide user high-level loop transformation directives for fusion and tiling.
- Developer describes loop bounds and data accesses.
- Compiler uses information to perform transformations legally.

Loop Chain Annotations

```
for( int i = lb; i <= ub; i += 1 )
  A[i] = B[i-1] + B[i] + B[i+1];</pre>
```

```
for( int i = lb; i <= ub; i += 1 )
    A[i] = A[i] / 3.0;</pre>
```

Loop Chain Annotations



Loop Chain Fusion

schedule(fuse())



Loop Chain Fusion with Explicit Shift

schedule(fuse((0), (1)))



Loop Chain Tiling

schedule(tile((2), parallel, serial))



After Tiling

Composition of Multiple Sequential Loop Optimizations

- How should sequential loop optimizations compose?
 - Previous loop chain tool easily composes optimizations.
 - How should we do this in OpenMP?
 - Unify into one framework?
 - Piecemeal?
 - Something in between?

Loop Chain Fuse + Tile

schedule(fuse(), tile((2), parallel, serial))



Loop Chain Tile + Fuse

schedule(tile((2), parallel, serial), fuse())



Interaction With Other OpenMP Constructs

How do sequential loop optimizations interact with different OpenMP constructs?

- For example: what would loop fusion on two OpenMP parallel loops do?
 - Union clauses?
 - \circ Invalid?
- Cross product of constructs is large.



- OpenMP API can provide methods of prescribing sequential loop optimizations.
- Sequential loop optimizations have positive performance impacts for parallel code.
- Developer can augment portions of analysis.
- API can provide multiple pathways for performing same transformations.

Discussion

- How could we compose different sequential loop optimizations?
 - Unified framework (à la loop chain)?
 - Piecemeal?
 - Something in between?
- How do we deal with interactions between different sequential loop and other OpenMP constructs?
 - Unified model?
 - Explicitly connect different components?
- What is the set of important sequential loop optimizations?