CSc 120 Introduction to Computer Programming II

Adapted from slides by Dr. Saumya Debray

10: Linked Lists

Python lists: reprise



concatenating two lists: O(n)

Question: Can we do insertion and concatenation in O(1) time? (complexity of other operations may change). ⇒ "Linked list"

Python lists: reprise

- Key feature: L[i] and L[i+1] are adjacent in memory
- This makes accessing L[i] very efficient
 - O(1)
- Insertion and concatenation require moving O(n) elements

 O(n)



- To get O(1) insertion and concatenation, we cannot afford to move O(n) list elements
- We have to relax the requirement that ith element is adjacent to (i+1)st element
 - any element can be anywhere in memory
- Each element has to tell us where to find the next element



• Linked list:

A collection of elements where each element has a value and a reference to the next element.

There is at least one variable that references the beginning of the list.

Each element of the list has a reference to the next list element



With each element of the list, keep a reference to the next list element



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Let's explore this idea using a file for a "node"

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Let's explore this idea using a file for a "node"

each file has two lines:

- the first line is a value
- the second line is a reference to the next "node" (file)

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 - each file has two lines:
 - the first line is a value
 - the second line is a reference to the next "node" (file)

Sample file "node": filename is 24.txt

value: aaa

next: 3.txt

EXERCISE

• Exploring linked lists using files as nodes

How would we add the word "total" to our linked lists of files so that the sentence reads:

The expert in anything was once a total beginner.

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• Exploring linked lists using files as nodes

How would we add the word "total" to our linked lists of files so that the sentence reads:

The expert in anything was once a total beginner.

Create a new "node" (a new file)

- The first line is "total"
- The second line is 19.txt
- What else do we have to do?
 - modify the file node for the value "a" to change its reference

References are addresses in memory.

Here is the diagram with explicit addresses (simplified).



Consider inserting a new node into the linked list



Specifically, add a new node between "bbb" and "ccc". What do we change?



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We want to add a new node between "bbb" and "ccc". What do we change?





Set the next references appropriately. What is the complexity of insertion?



*assuming we have a reference to the node of insertion



To insert an element into a linked list: set next references appropriately



Concatenation



To concatenate two linked lists: set next reference of end of first list to refer to beginning of second list



* once we have a reference to the end of the first list

implementation

Nodes: Implementation

class Node:

def __init__(self, value):
 self._value = value # reference to the object at that node
 self._next = None # reference to the next node in the list

Getters:

```
def value(self):
    return self._value
```

Setters:

```
def set_value(self, value):
    self._value = value
```

```
def next(self):
return self._next
```

```
def set_next(self, next):
    self._next = next
```

A linked list is just (a reference to) a sequence of nodes



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class LinkedList:

def __init__(self):
 self._head = None

def is_empty(self):
 return self._head == None

def head(self): return self._head

addition at the head of the list

Adding a node at the head



Adding a node at the head



Sequence of operations for an add method:

1. new._next = L._head



2. L._head = new



Adding a node at the head

class LinkedList:

def __init__(self):
 self._head = None

O(1)

add a node new at the head of the linked list
def add(self, new):
 new._next = self._head
 self. head = new

```
class Node:
    def __init__(self, value):
        self._value = value
        self._next = None
    ...
class LinkedList:
```

```
def __init__(self):
    self._head = None
```

```
def add(self, new):
    new._next = self._head
    self._head = new
```

infile = open("infile.txt")
my_list = LinkedList()
for line in infile:
 this_node = Node(line)
 my_list.add(this_node)

,	<u>infile.txt</u>
	аа
	bb
	СС

```
class Node:
    def __init__(self, value):
        self._value = value
        self._next = None
    ...
class LinkedList:
    def __init__(self):
        self._head = None
    def add(self, new):
        new._next = self._head
        self._head = new
```

```
infile = open("infile.txt")
my_list = LinkedList()
for line in infile:
    this_node = Node(line)
    my_list.add(this_node)
```

my_lis <mark>t</mark>	
	LinkedList
_head	None

infile.txt		
аа		
bb		
СС		


























Adding a node at the head

Changing the order of assignments does not work:

```
def broken_add(self, new):
```

self._head = new

new._next = self._head

def add(self, new):

new._next = self._head

self._head = new



appending to the tail of the list

Adding a node at the tail

To add a node X at the end (i.e., tail) of a list L:

1. find the last element Y of L

2. Y._next = X

Adding a node at the tail

To add a node X at the end (i.e., tail) of a list L:

- 1. find the last element Y of L
- 2. Y.__next = X _____

O(n)

O(1)

Adding a node at the tail

To add a node X at the end (i.e., tail) of a list L:

- 1. find the last element Y of L
- 2. Y._next = X

Gotchas to watch out for:

- what if there is no last element?
 - how can we tell?
 - what should we do?



EXERCISE

- Consider a linked list whose value attributes consist of strings.
- Write a method replace (arg1, arg2) that replaces the value attributes of all nodes that equal arg1 with arg2.

finding the nth element

Finding the nth element

class LinkedList:

```
# return the node at position n of the linked list
def get_element(self, n):
    elt = self._head
    while elt != None and n > 0:
        elt = elt._next
        n -= 1
    return elt
```

insertion

Suppose we want to insert a node X into a list here:



Then we have to adjust the next-node reference on the node Y just before that position

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The order of operations is important:



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- 1. X._next = Y._next
- 2. Y.__next = X

Inserting a node X at position *n* in a list L:

- find the node Y at position *n*−1
 - iterate n-1 positions
 from the head of the list*
- 2. insert X after Y
 - adjust next-node references as in previous example

 * do something sensible if the list has fewer than n-1 nodes $Y = L._head$ for i in range(n-1): O(n) $Y = Y._next$ X._next = Y._next O(1)

$$Y._next = X$$

class LinkedList:

```
# insert a node new at position n
def insert(self, new, n):
    if n == 0:
        self.add(new)
    else:
        prev = self.get_element(n-1)
        new.next = prev.next
        prev.next = new
```

deletion



Suppose we want to delete this node:





Deleting a node

Suppose we want to delete this node:



- 1. find the node Y just before X
 (i.e., Y._next == X)
- 2. Y._next = X._next
- 3. X._next = None

O(n)

O(1)

Deleting a node

```
class LinkedList:
```

```
# delete a node X
def delete(self, X):
   if self. head == X:
                            # X is the head of the list
      self. head = X. next
   else:
      Y = self. head
      while Y._next != X:
          Y = Y. next
      Y. next = X. next
   X.next = None
```

deletion (revisited)



Suppose we want to delete this node:





Suppose we want to delete this node:





def delete(self, x):
 r = self._head
 while r != None:
 if r == x:
 <delete node x>
 return
 r = r. next

- Does this code pattern work for delete?
- It worked for len, replace, count_vowels ...



r = self._head while r != None: if r == x: <delete node x> return r = r._next

- No, does not work
- We need a reference to the previous node

Deleting a node

```
class LinkedList:
```

```
# delete a node X
def delete(self, X):
   if self. head == X:
                            # X is the head of the list
      self. head = X. next
   else:
      Y = self. head
      while Y._next != X:
          Y = Y. next
      Y. next = X. next
   X.next = None
```
concatenation

Concatenating two linked lists

class LinkedList:

concatenate list2 at the end of the list def concat(self, list2): if self. head == None: *# list is empty* self. head = list2. head else: tail = self._head while tail. next != None: O(n) O(1) tail = tail. next tail.next = list2. head

maintaining a tail reference

Maintaining a tail reference

A variation is to also maintain a reference to the tail of the list













Maintaining a tail reference

```
• Concatenation and append become O(1):
def concat(self, list2):
    if self. head == None:
       self. head = list2. head
       self. tail = list2. tail
    else:
        self. tail. next = list2. head
        self. tail = list2. tail
```

• All linked list operations must now make sure that the tail reference is kept properly updated

Linked lists: summary

Operation	Without tail reference	With tail reference
add to front of list	O(1)	
append to end of list	O(n)	O(1)
find nth element	O(n)	
insert	O(1) if prev. node is available O(n) otherwise	
delete	O(1) if prev. node is available O(n) otherwise	
concatenate	O(n)	O(1)