

## Maps:

- We know an array allows us to use an index to access a value in constant time.
- However, arrays don't do us any good if we don't have an index.
- A map lets us take a key and access a value.
  - The key doesn't have to be an index. It could be a string as in the happiness program.
- How long does it take to look up a value given a key?
  - It depends on how we implement the map.
    - Suppose we implement the map with a linked list.
    - Can we do this?
      - Yes
    - How long does it take to get the value given a key?
      - It is an  $O(n)$  operation where  $n$  is the number of items stored

## Maps:

- Now suppose we implement the map with a Binary Search Tree.
  - How long does it take to look up a value?
    - It is  $O(\log n)$  if the tree is what?
      - balanced
  - Binary Search Trees are actually quite fast but
    - It requires some work to keep them balanced
    - $O(\log n)$  is fast, but not as fast as  $O(1)$  (constant time)

## Hash Tables:

- A *hash table* or *hash map* uses the keys to sort the values into *buckets*.
- For example, we might sort your test papers into separate piles based on the first letter of your last name.
  - Then to get your test paper I could pick up the correct pile (constant time) and search through only a subset of the class.
  - How much time do you expect it to take? (169 students)
    - $169/26 = 6.5$  (expect  $\frac{1}{2}$  piles to have 6 and  $\frac{1}{2}$  to have 7)
    - Actually, the time would vary depending on the letter

A	5
B	8
C	7
D	8
E	1
F	5
G	12
H	8
I	1
J	2
K	5

L	19
M	11
N	5
O	2
P	10
Q	1
R	12
S	11
T	5
U	1
V	6

W	11
X	2
Y	4
Z	7

These are the actual numbers of students whose names start with these letters. Notice that they are not evenly distributed. The E pile only contains 1, but the L pile has 19.

## Hash Tables:

- A *hash function* is a function that takes a key and returns the index of a bucket.
  - In the previous example the function takes the first letter of the name and returns a value between 0 and 25.
- Ideally we would like the hash function to spread the keys out evenly over the buckets.

## Hash Tables:

- Let's suppose our keys are strings, what are some possible hash functions?
  - Length of string.
    - This will have a lot of clustering. All strings of the same length map to the same index.
  - First letter of string. - we've already looked at this
  - Sum of Unicode values of the characters.
    - This is better than the other two. Sill, all anagrams map to the same index.
    - Note also that this function can give us very large numbers. To convert this number to a valid index, we can mod it by the number of buckets.
  - We could multiplying the char value times it's position.
- In general a good hash function tries to use all the information in the key.

## Hash Tables:

- A hash map has an array of buckets.
- Each bucket is a simple structure usually a linked list.
  - Even the best hash map could sometimes map different keys to the same index, so you need to represent the bucket with something like a linked list that can support multiple entries.
- A hash map also has a hash function which maps a key to an index.

## Hash Tables:

- The functions of a Hash Map work like:

**hashPut (key, value)**

**i = hashMap(key)**

**add node(key, value) to linked list to bucket i**

**hashGet (key)**

**i = hashMap(key)**

**for each node n in linked list of bucket i**

**if n.key = key**

**return value**

**return value not found**

## Hash Tables:

- How many buckets should you have?
- Ideally we would like a Hash Map to work in constant time, so we'd like one bucket per item stored.
  - Notice that if each bucket has only one item in it, then the search time will be constant.
- Some implementations of hash maps resize the structure automatically when too many items are stored.
- How could we implement that?