

# Algorithms CSs545 — Homework #6

Due: 12/11/02.

December 10, 2002

1. CLR 26-1. (CLRS 25-1).
2. The question refers to the stable matching algorithm as studied in class. We define a pair  $(m_i, w_j)$  to be a *stable pair* if there exists at least one stable marriage at  $m_i$  and  $w_j$  are matched. Prove that during the algorithm studied in class, if  $w_j$  rejects  $m_i$ , then they are not a stable pair. Hint — assume this is not the case, consider the first event at which a woman  $w_1$  rejects a man  $m_1$  while  $(m_1, w_1)$  are stable pair. Let  $m_2$  be the man whose proposal to  $w_1$  caused her to reject  $m_1$ , and show contradiction.
3. Assume that the coin used for setting the level in the SkipList data structure is not biased, and fall in head with probability  $\alpha$ , which is not necessarily  $1/2$ . What is the time for inserting an element - and the number of levels it participates in ? What is the effect on the query time, and on if  $\alpha$  is much larger than  $1/2$  ? What is the effect if it is much smaller ?
4. Explain how would you use SkipList to store a set  $S$  of numbers, in a way that would allow you to find the number of elements in  $S$  which are smaller than a query value  $x$ . You should be able to add elements to  $S$ , delete elements, and perform a query in time  $O(\log |S|)$ .
5. Let  $G(V, E)$  denote a graph with weights assigned to its edges, where  $V = \{v_1 \dots v_n\}$ . Read and understand how we can find a matrix  $\Pi[\cdot, \cdot]$ , such that  $\Pi[i, j]$  contains the vertex preceding  $v_i$  on the shortest path from  $v_i$  to  $v_j$ .
6. CLR 27.2-4 (CLRS 26.2-4)
7. CLR 27.2-9 (CLRS 26.2-)
8. (only if you feel like it) Let  $L$  be singly connected linked list.  $L$  is called a *snake* if the last element of  $L$  points to NULL, and is called a *snail* if the last element points to one of the last elements of  $L$ . You are given a pointer to the first

element of  $L$ , and  $O(1)$  additional memory. You cannot change  $L$  itself - not even temporarily. Find in  $O(n)$  time (a) if  $L$  is a snake or a snail, and (b) how many elements are there in  $L$ . There is a solution that does not use unbounded search.