Why types?

Types save typing.

What does `a+b` mean?

In Java it could be
1. `a +int b`.
2. `a +float b`.
3. `a concat_string b`.
4. `int2float(a) +float b`.
5. `a +float int2float(b)`.
6. `int2string(a) concat_string b`.

etc, all depending on the types of `a` and `b`.

Why types...?

In Icon variables are not given explicit types. Instead, operations carry the types:
1. `a | b` means binary or on integers.
2. `a || b` means string concatenation.
3. `a ||| b` means list concatenation.

Icon has lots of operators...

In other words, without types, we would have to be much more explicit about which operations are performed where.

Why types...?

Icon programs become a bit wordier since every operator effectively encode the required type of the operands.

On the other hand, it also becomes more readable since we can see directly from the operator what operation will be performed.

```
global x,y,z
procedure p()
x := x + y  # integer addition
x := x || y  # string concatenation
x := x ||| y # list concatenation
end
```
Why types...?

To figure out which operation is performed in a Java program, we have to find the declarations of all variables to find their declared type:

```java
int x;
String y;
float z;
void p() {
    x = x + 5; /* integer addition */
    z = z + 5.0; /* float addition */
    y = y + "X"; /* string concatenation */
}
```

Types prevent errors. Types save the programmer from himself. Types prevent us from adding a character and a record.

```java
int A[20];
float x;
void p() {
    A[x] = 5;
    x = x + A;
}
```

Why types...?

Types permit optimization. A compiler can generate better code for `a+b` if it knows that both variables must be integers, than if the exact types aren’t known until runtime:

```java
global a,b
procedure p() {
    a = new array [20]
    ...
    b = new array [20]
    ...
    a = a + b /* what operation is performed */
}
```

Type Systems

A type system consists of:
- a mechanism for defining types,
- rules for type equivalence,
- rules for type compatibility,
- rules for type inference.
Type Systems...

*Type equivalence* determines when the types of two values are the same:

```
TYPE A = ARRAY [0..10] OF CHAR;
TYPE B = ARRAY [0..10] OF CHAR;
VAR a : A;
VAR b : B;
BEGIN
  a := b; (* legal? *)
END
```

Are the types of `a` and `b` the same?

Type Systems...

*Type compatibility* determines when a value of a given type can be used in a given context:

```
VAR a : float;
VAR b : int;
BEGIN
  a := a + b;
END
```

Can you add an `int` and a `float`?

Type Systems...

*Type inference* defines the type of an expression based on its parts and surrounding context:

```
global a, b, c
procedure p(x)
  if x = 5 then
    a := x
  else
    a := "hello"
  write(a)
end
procedure main()
p(5)
end
```

What type of data can be written here?

Type Checking

*Type checking* ensures that a program obeys a language’s type rules.

A *type clash* is a violation of the typing rules.

```
class C {
  void p() {
    int x = new C();
  }
}
```
Type Checking — Strong Typing

A language is **strongly typed** if

- ⊕ is an operator in \( L \) that expects an object of type \( T \),
- \( L \) prohibits ⊕ from accepting objects of any other type,
- and \( L \) requires an implementation (a compiler, interpreter, etc) to enforce this prohibition.

In other words, a strongly typed language does not allow us to perform operations on the “wrong” type of data.

Type Checking — Weak Typing

In a **weakly typed** language there are ways to “escape” the type system.

In C, for example, it is possible to cast a pointer to a float, add 3.14 to it, and cast it back to a pointer:

```c
int main() {
    int* p = (int*) malloc (sizeof(int));
    float f = *((float*) &p) + 3.14;
    p = (int*)(*((int*) &f));
}
```

Such operations are probably meaningless and a strongly typed language would prohibit them.

Type Checking — Static/Dynamic Typing

- A language is **statically typed** if type checking is done at compile-time.
- A language is **dynamically typed** if type checking is done at run-time.
- In practice, even languages which are considered statically typed do some checking at run-time.
- Languages can usually be classified as **mostly strongly typed**, **mostly statically typed**, etc.

Terminology

Benjamin C. Pierce has said:

I spent a few weeks . . . trying to sort out the terminology of **strongly typed, statically typed, safe**, etc., and found it amazingly difficult. . . . The usage of these terms is so various as to render them almost useless.

It is possible to say

My language is more strongly typed than your language.

but harder to argue that

My language is strongly typed/statically typed, etc.
Examples — Pascal

- Pascal is mostly strongly and statically typed.
- **Untagged variant records** are a loophole. They allow us to turn a value of one type into an object of some unrelated type.
- Unlike C, array bounds are checked.

Pascal – Untagged Variant Records

```pascal
type rec = record
  a : integer;
  case boolean of
    true : (x : integer);
    false : (y : char);
  end;
end;

var r: rec;
begin
  r.x := 55; r.y := 'A'; write(r.x);
end.
```

- This construct is used to bypass Pascal’s strong typing.

Examples — C

- C is weakly and statically typed.
- Pointers can be cast willy-nilly which makes it easy to bypass the type system.
- Array references are not checked:
  ```c
  int main() {
    int A[20];
    int B[20];
    A[25] = 5;
  }
  
  Negative indices were used in the old days to overwrite the operating systems.
  
  Today, buffer overflows are how most viruses compromise security.
  ```

Examples — Ada

- Ada is strongly and mostly statically typed.
- Unlike Pascal, variant records must be tagged:
  ```ada
  type Device is (Printer, Disk, Drum);
  type Peripheral(Unit : Device := Disk) is record
    case Unit is
      when Printer => Line_Count : Integer;
      when others => Cylinder : CIndex;
    end case;
  end record;
  ```

- Today, buffer overflows are how most viruses compromise security.
Examples — Ada...

- It is, however, possible to do **non-converting casts** (similar to C), but in a very explicit way:

  ```
  function float2int is 
      new unchecked_conversion(float, integer);
  ... 
  f := float2int(i);
  ```

- Some errors can’t be checked at compile-time:

  ```
  I, J : Integer range 1 .. 10 := 5; 
  K : Integer range 1 .. 20 := 15; 
  I := J; -- identical ranges 
  K := J; -- compatible ranges 
  J := K; -- will raise an exception if K>10
  ```

Examples — Scheme

- Scheme is completely dynamically typed, so programmers often insert extra checks:

  ```
  (define (sum l) 
    (cond 
      ((null? l) 0) 
      ((not (list? l)) 
        (error "list expected")) 
      ((not (number? (car l))) 
        (error "list of numbers expected") 
      (else (+ (car l) (sum (cdr l)))) 
    ))
  ```

Examples — Java

- Java is strongly and mostly statically typed.

- An exception is thrown here because an A-object can’t be cast to a B-object:

  ```
  class A {} 
  class B extends A { 
      int x; 
  } 
  void p() { 
      B b = (B) new A(); 
  }
  ```

Typing

```
<table>
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<tr>
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<th>static</th>
<th>dynamic</th>
<th>strong</th>
<th>weak</th>
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</tbody>
</table>
```
Type Inference

In statically typed languages types are inferred in the compiler, before the program is run:

```pascal
procedure p (x : integer);
var z : real;
var c : char;
begin
  write(x + z); /* convert x to real, write a real */
  write(c + z); /* type error */
end
```

Haskell and similar languages don’t require the programmer to give types to variables and functions. Instead, the compiler infers types. Given

```haskell
len [] = 0
len _:xs = 1 + len xs
```

the Haskell translator will infer a most general type:

```haskell
len :: [a] -> Int
```

Haskell is strongly and statically typed, although the programmer rarely have to provide explicit type information.

So, What is a Type?

There are three ways to think about types:

1. **denotational view** — a type is a set of values;
2. **constructive view** — a type is what we can construct from the type constructors in the language;
3. **abstraction-based view** — a type denotes a data object and a well-defined set of allowable operators on this object.

At different times, we may look at a type in any of these ways.

Denotational View

A type $T$ is a set of values $\{t_0, t_1, t_2 \ldots \}$. A value $v$ is of type $T$ if it belongs to the set. A variable $v$ is of type $T$ if it is guaranteed to always hold a value in the set.

A `char` type in Pascal is the set of 128 seven-bit ASCII characters:

$$
\{\ldots,
  "0", \ldots, "9", \ldots,
  "A", \ldots, "Z", \ldots,
  "a", \ldots, "z", \ldots\}
$$
Constructive View

A Pascal type is (roughly)

\[ \text{type ::= } \]
\[ \text{integer | real | char | boolean ... } \]
\[ \text{[ expr .. expr ] | } \]
\[ \text{SET OF type | } \]
\[ \text{ARRAY type OF type | } \]
\[ \text{RECORD [field_list] END} \]

I.e., a Pascal type is either one of the built-in types, or ones we define ourselves by composing type constructors, such as ARRAY, RECORD, etc:

\[ \text{END T = RECORD} \]
\[ \text{a : real; } \]
\[ \text{b : ARRAY ["a".."z"] OF SET OF char;} \]
\[ \text{END;} \]

Abstraction-Based View

A type is an abstract data type.

The next slides shows what the Modula-3 language manual says about the operations that are allowed on Words.

The allowed operations include arithmetic and logical operations.

There is no “pointer dereferencing” operation defined, however, so apparently this operation is not allowed.

Abstraction-Based View...

INTERFACE Word;

TYPE T = INTEGER;

PROCEDURE Plus (x,y: T): T;
PROCEDURE Times (x,y: T): T;
PROCEDURE Minus (x,y: T): T;
PROCEDURE Divide(x,y: T): T;
PROCEDURE Mod(x,y: T): T;
PROCEDURE LT(x,y: T): BOOLEAN;
PROCEDURE LE(x,y: T): BOOLEAN;
PROCEDURE GT(x,y: T): BOOLEAN;
PROCEDURE GE(x,y: T): BOOLEAN;
PROCEDURE And(x,y: T): T;
PROCEDURE Or (x,y: T): T;
PROCEDURE Xor(x,y: T): T;
PROCEDURE Not (x: T): T;
PROCEDURE Shift(x: T; n: INTEGER): T;
PROCEDURE Rotate(x: T; n: INTEGER): T;
PROCEDURE Extract (x: T; i, n: CARDINAL): T;
PROCEDURE Insert(x: T; y: T; i, n: CARDINAL): T;

END Word.

Readings and References

Read Scott, pp.307–312.