

CSc 520

Principles of Programming Languages

30 : Exceptions

Christian Collberg

collberg+520@gmail.com

Department of Computer Science
University of Arizona

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Exception Handling

- What should a program do if it tries to pop an element off an empty stack, or divides by 0, or indexes outside an array, or produces an arithmetic error, such as overflow?
- In C, many procedures will return a *status code*. In most cases programmers will “forget” to check this status flag.
- Modern languages have built-in *exception* handling mechanisms. When an exception is *raised* (or *thrown*) it must be handled or the program will terminate.
- Exceptions can be raised implicitly by the run-time system (overflow, array bounds errors, etc), or explicitly by the programmer.

Exception Handling...

- When an exception is raised, the run-time system has to look for the corresponding *handler*, the piece of code that should be executed for the particular exception.
- The right handler cannot be determined statically (at compile-time). Rather, we have to do a dynamic (run-time) lookup when the exception is raised.
- In most languages, you start looking in the current block (or procedure). If it contains no appropriate handler, you return from the current routine and re-raise the exception in the caller. This continues until a handler is found or until we get to the main program (in which case the program terminates with an error).

Exceptions in PL/I

So, what happens afterwards?

- What happens after an exception handler has been found and executed?
resumption model Go back to where the exception was raised and re-execute the statement (PL/I).
termination model Return from the procedure (or unit) containing the handler (Ada).

Exceptions in PL/I

- Exceptions are declared like this:

```
ON condition  
    statement
```

This statement is not actually executed. It's just “remembered” until later, should an error occur.
- For example:

```
ON OVERFLOW  
    GOTO error;  
  
...  
error:  
    PRINT "something bad happened"
```

Exceptions in PL/I...

- Consider this example:

```
...  
ON OVERFLOW  
    PRINT "use smaller numbers, please!"  
...  
A := A*10000000000000;
```
- Where does execution continue after the exception handler has executed? *After* the location where the exception was thrown.

Exceptions in Modula-3

Exceptions in Modula-3

- Exceptions are declared like this:

```
INTERFACE M;
  EXCEPTION Error(TEXT);
  PROCEDURE P () RAISES {Error};
END M;
```

- Exceptions can take parameters. In this case, the parameter to `Error` is a string. Presumably, the programmer will return the kind of error in this string.
- The declaration of `P` states that it can only raise one exception, `Error`.
- If there is no `RAISES` clause, the procedure is expected to raise no exceptions.

Exceptions in Modula-3...

- S_1 and S_2 can raise exceptions implicitly, or the programmer can raise an exception explicitly using `RAISE`.
- When the `Error`-exception is raised, the `EXCEPT`-block is searched and the code for the `Error` exception is executed.

```
PROCEDURE P () RAISES {Error};
BEGIN
  TRY
     $S_1$ ; RAISE Error("Help!");  $S_2$ ;
  EXCEPT
    Error (V) => Write(V); |
    Problem (V) => Write("No Probs!"); |
    ELSE Write("Unhandled Exception!");
  END;
END P;
```

Exceptions in Modula-3...

- An unhandled exception is re-raised in the next dynamically enclosing `TRY`-block. If no matching handler is found the program is terminated.

```
MODULE M;
BEGIN
  TRY
    TRY  $S_1$ ; EXCEPT
      Problem (V)=>Write(V);
    END;
  EXCEPT
    Error (V) => Write(V); |
    ELSE Write("Unhandled Exception!");
  END;
END M;
```

Exceptions in Modula-3...

- An unhandled exception is re-raised in the calling procedure. Exception handlers can explicitly re-raise an exception, or raise another exception.

```
MODULE M;
  PROCEDURE P ();
  BEGIN
    TRY  $S_1$ ; EXCEPT
      Problem (V)=>RAISE Error("OK")
    ENDEND P;
  BEGIN
    TRY P(); EXCEPT
      Error (V) => Write(V); |
      Problem (V) => Write(V);
    END;
  END M;
```

Exceptions in Java

- In Java, exceptions are classes. They are declared like this:

```
class StupidError extends Exception {}
```

- An exception is thrown by creating a new exception object, and then calling `throw`:

```
throw new StupidError("Forgot to buy milk")
```

Exception hierarchy

- Java defines a hierarchy of exceptions. Programmers can construct their own by extending one of these classes:

```
class Throwable {}
class Exception extends Throwable {}
class InterruptedException extends Exception {}
class RuntimeException extends Exception {}
class ArithmeticException extends RuntimeException {}
class NullPointerException extends RuntimeException {}
class ClassCastException extends RuntimeException {}
class Error extends Throwable {}
class ThreadDeath extends Error {}
```

Catching an exception

- Exceptions are caught like this:

```
try {
    throw new Exception("They were out of milk")
} catch (StupidError e) {
    ...
} catch (Exception e) {
    ...
} catch (Throwable e) {
    ...
}
```

- Each `catch`-clause is considered in turn until one is found that is a subclass of the exception object thrown.

The finally-clause

- The `finally`-clause is executed regardless of whether an exception is thrown or not. This can be used to close files, etc:

```
try {  
    throw new Exception("They were out of mil  
} catch (StupidError e) {  
    getACow();  
} finally {  
    eatCornFlakes();  
}
```

Implementing exceptions

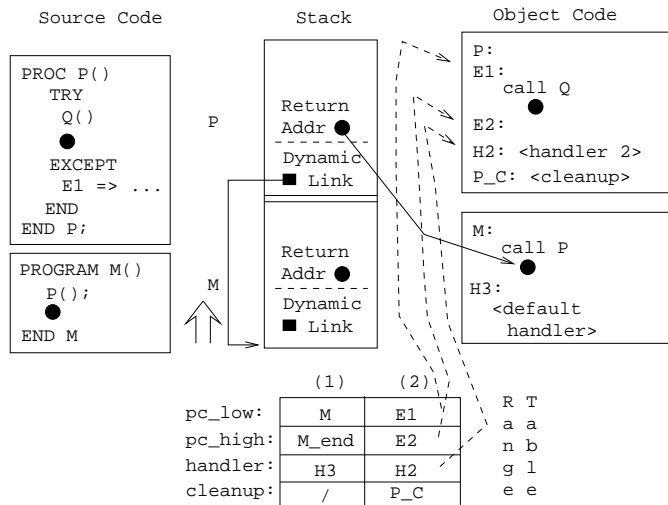
Implementation

- We want 0-overhead exception handling. This means that – unless an exception is raised – there should be no cost associated with the exception handling mechanism.
- We allow raising and handling an exception to be quite slow.
- When an exception is raised we need to be able to
 1. in the current procedure find the exception handler (if any) that encloses the statement that raised the exception, and
 2. rewind the stack (pop activation records) until a procedure with an exception handler is found.

The Range Table

- We build a *RangeTable* at compile-time. It has one entry for each procedure and for each **TRY**-block.
- Each entry holds four addresses: `pc_high`, `pc_low`, `handler` and `cleanup`.
- `[pc_low...pc_high]` is the range of addresses for which `handler` is the exception handler.

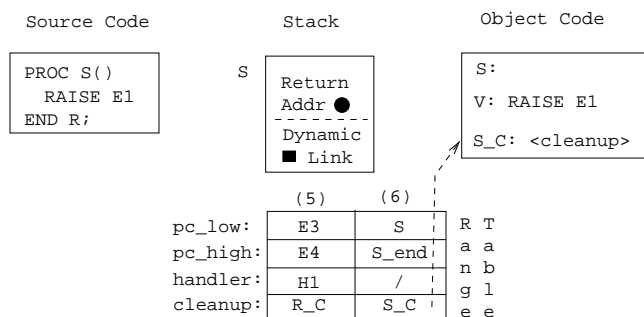
The Range Table...



Unwinding the Stack(Locate)

- Let procedure S raise exception E at code address V . We search the range table to find an entry which covers V , i.e. for which $pc_low \leq V \leq pc_high$.
- Entry (6) covers all of procedure S (for S to S_end), and hence V . There's no exception handler for this range. We just execute S 's cleanup code, S_C .
- S_C will restore saved registers, etc, and deallocate the activation record.

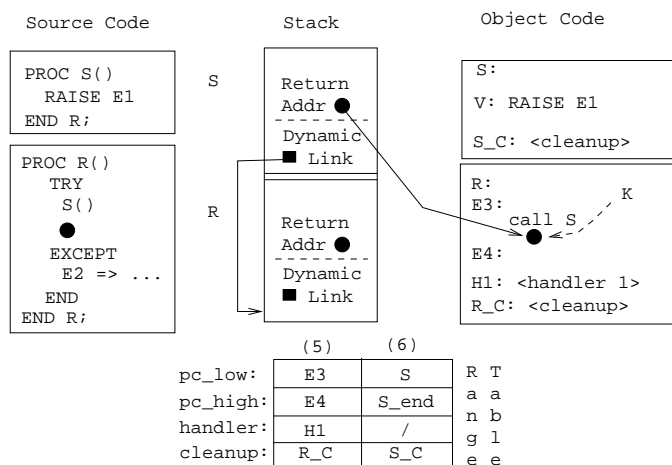
Unwinding the Stack (Locate)...



Unwinding the Stack (Unwind)

- Since S didn't have a handler, we must unwind the stack until one is found.
- S 's return address is K , which is covered by entry (5) in the range table. Entry (5) has a handler defined (at address $H1$). Run it!

Unwinding the Stack (Unwind)...



The Exception Handler

- The exception handler itself can be translated as a sequential search.
- If the **TRY-EXCEPT**-block has no **ELSE** part, the default action will be to re-raise the exception.

TRY	$S_1;$	$S_1;$
	$S_1;$	RAISE e ;
	RAISE e ;	$S_2;$
	$S_2;$	\Rightarrow
EXCEPT		IF $e = E_1$ THEN H_1
	$E_1 \Rightarrow H_1$	ELSIF $e = E_2$ THEN H_2
	$E_2 \Rightarrow H_2$	ELSE
END;		RAISE e
		ENDIF;

The Algorithm

```

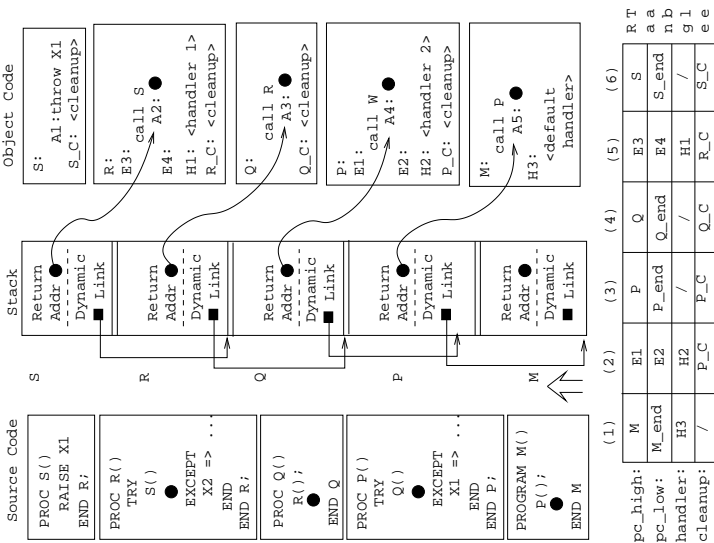
LOOP
  D := The first procedure descriptor (Range Table
        entry) such that D.pc.low <= PC <= D.pc.high;
  IF D.handler = the default handler THEN
    abort and coredump
  ELSIF D.handler ≠ NIL THEN GOTO D.handler;
  ELSE
    Execute the cleanup routine D.cleanup;
    PC := Return address stored in the current frame;
    SP := SP of previous frame;
    FP := FP of previous frame;
  END;
END;

```

Example — Explanation of source code

- Consider the example on the next slide.
- The main program calls procedure $P()$. There is a <default handler> defined for the program at address H3.
- Procedure $P()$ calls $Q()$. Exception X1 is caught by the handler at address H2.
- $Q()$ calls $R()$.
- $R()$ calls $S()$. Exception X2 is caught by the handler at address H1.
- $S()$ throws exception X1 at address A1.

Example



Example — Explanation of Actions

- $A1 \in [S, S_end]$, in Range Table entry (6). (6) has no handler, so we execute its cleanup routine (S_C) and update PC to the return address, $A2$.
- Since $A2 \in [E3, E4]$ in Range Table entry (5), and $(5).handler == H1 \neq \text{NIL}$, we GOTO $H1$. This handler doesn't handle exception $X1$, so it will simply re-raise $X1$.
- $Q()$ has no handler, so we execute its cleanup routine (Q_C) and propagate the exception to $P()$. I.e. We update PC to the return address stored in Q 's frame, $A4$.
- Since $A4 \in [E1, E2]$ in Range Table entry (2), and $(2).handler = H2$, we GOTO $H2$. This handler catches $X1$. \Rightarrow Done.

Exceptions in C

- In C, `setjmp/longjmp` can be used to implement exceptional control flow:

```
if (!setjmp(buffer)) {
    /* setjmp returned 0. Protected code.*/
    ...
    longjmp(buffer);
    ...
} else {
    /* setjmp returned 1. Handler code. */
}
```


- The first time `setjmp` returns 0 and execution continues as normal. When `longjmp` is called it appears as if `setjmp` has returned for the second time, this time returning 1. The state is now the same as it was when `setjmp` was first called.
- `setjmp`'s buffer argument stores the program's current state, in particular register values.
- Unlike a “real” exception handler, the stack is not rewind nicely. Rather, all stack frames are thrown away. This can lead to problems if not all register values have been saved back in memory. Variables that may be thus affected should be declared `volatile`, i.e. they will always be returned to memory after operated on.

- Read Scott: pp. 441–452
- Drew, Gough, Lederman, *Implementing Zero Overhead Exception Handling*,
<http://www.dstc.qut.edu.au/~gough/zeroex.ps>.
- Drew, Gough, *Exception handling: Expecting the Unexpected*, Computer Language, Vol 32, No 8, pp. 69–87, 1994.

Summary

- The algorithm we've shown has no overhead (not even one instruction), unless an exception is thrown.
- The major problem that we need to solve is finding the procedure descriptor for a particular stack frame.
- An alternative implementation would be to store a pointer in each frame to the appropriate descriptor. The extra space is negligible, but it would cost 1-2 extra instructions per procedure call.