CSc 520

Principles of Programming Languages

30 : Exceptions

Christian Collberg

collberg+520@gmail.com

Department of Computer Science

University of Arizona

Exception Handling

- What should a program do if it tries to pop an element off an empty stack, or divides by 0, or indexes outside an array, or produces an arithmetic error, such as overflow?
- In C, many procedures will return a *status code*. In most cases programmers will “forget” to check this status flag.
- Modern languages have built-in *exception* handling mechanisms. When an exception is *raised* (or *thrown*) it must be handled or the program will terminate.
- Exceptions can be raised implicitly by the run-time system (overflow, array bounds errors, etc), or explicitly by the programmer.

Exception Handling. . .

- When an exception is raised, the run-time system has to look for the corresponding *handler*, the piece of code that should be executed for the particular exception.
- The right handler cannot be determined statically (at compile-time). Rather, we have to do a dynamic (run-time) lookup when the exception is raised.
- In most languages, you start looking in the current block (or procedure). If it contains no appropriate handler, you return from the current routine and re-raise the exception in the caller. This continues until a handler is found or until we get to the main program (in which case the program terminates with an error).

Exceptions in PL/I

So, what happens afterwards?

- What happens after an exception handler has been found and executed?

  resumption model Go back to where the exception was raised and re-execute the statement (PL/I).

  termination model Return from the procedure (or unit) containing the handler (Ada).

Exceptions in PL/I

- Exceptions are declared like this:

  ```
  ON condition
     statement
  ```

  This statement is not actually executed. It's just "remembered" until later, should an error occur.

- For example:

  ```
  ON OVERFLOW
     GOTO error;
  ...
  error:
     PRINT "something bad happened"
  ```

Exceptions in PL/I. . .

- Consider this example:

```
...
ON OVERFLOW
   PRINT "use smaller numbers, please!"
...
A := A*1000000000000;
```

- Where does execution continue after the exception handler has executed? *After* the location where the exception was thrown.

Exceptions in Modula-3

Exceptions in Modula-3

- Exceptions are declared like this:

```
INTERFACE M;
   EXCEPTION Error(TEXT);
   PROCEDURE P () RAISES {Error};
END M;
```

- Exceptions can take parameters. In this case, the parameter to `Error` is a string. Presumably, the programmer will return the kind of error in this string.
- The declaration of `P` states that it can only raise one exception, `Error`.
- If there is no `RAISES` clause, the procedure is expected to raise no exceptions.

Exceptions in Modula-3...

- S_1 and S_2 can raise exceptions implicitly, or the programmer can raise an exception explicitly using RAISE.
- When the Error-exception is raised, the EXCEPT-block is searched and the code for the Error exception is executed.

```
PROCEDURE P () RAISES {Error};
BEGIN
  TRY
    S1; RAISE Error("Help!"); S2;
  EXCEPT
    Error (V) => Write(V); |
    Problem (V) => Write("No Probs!"); |
    ELSE Write("Unhandled Exception!");
  END;
END P;
```

Exceptions in Modula-3. . .

- An unhandled exception is re-raised in the next dynamically enclosing TRY-block. If no matching handler is found the program is terminated.

```
MODULE M;
BEGIN
   TRY
      TRY S1; EXCEPT
         Problem (V)=>Write(V);
      END;
   EXCEPT
      Error (V) => Write(V); |
      ELSE Write("Unhandled Exception!");
   END;
END M;
```

Exceptions in Modula-3. . .

- An unhandled exception is re-raised in the calling procedure. Exception handlers can explicitly re-raise an exception, or raise another exception.

```
MODULE M;
   PROCEDURE P ();
   BEGIN
      TRY S1; EXCEPT
         Problem (V)=>RAISE Error("OK")
      END;END P;
BEGIN
   TRY P(); EXCEPT
      Error (V) => Write(V); |
      Problem (V) => Write(V);
   END;
END M;
```

Exceptions in Java

Exceptions as classes

- In Java, exceptions are classes. They are declared like this:

  ```
  class StuipdError extends Exception {}
  ```

- An exception is thrown by creating a new exception object, and then calling `throw`:

  ```
      throw new StupidError("Forgot to buy milk
  ```

Exception hierarchy

- Java defines a hierarchy of exceptions. Programmers can construct their own by extending one of these classes:

```
class Throwable {}
   class Exception extends Throwable {}
      class InterruptedException extends Exception {}
      class RuntimeException extends Exception {}
         class ArithmeticException extends RuntimeExceptio
         class NullPointerException extends RuntimeExcepti
         class ClassCastException extends RuntimeException
   class Error extends Throwable {}
      class ThreadDeath extends Error {}
```

Catching an exception

- Exceptions are caught like this:

```
try {
   throw new Exception("They were out of mil
} catch (StupidError e) {
   ...
} catch (Exception e) {
   ...
} catch (Throwable e) {
   ...
}
```

- Each `catch`-clause is considered in turn until one is found that is a subclass of the exception object thrown.

The `finally`-clause

- The `finally`-clause is executed regardless of whether an exception is thrown or not. This can be used to close files, etc:

```
try {
   throw new Exception("They were out of mil
} catch (StupidError e) {
   getACow();
} finally {
   eatCornFlakes();
}
```

Implementing exceptions

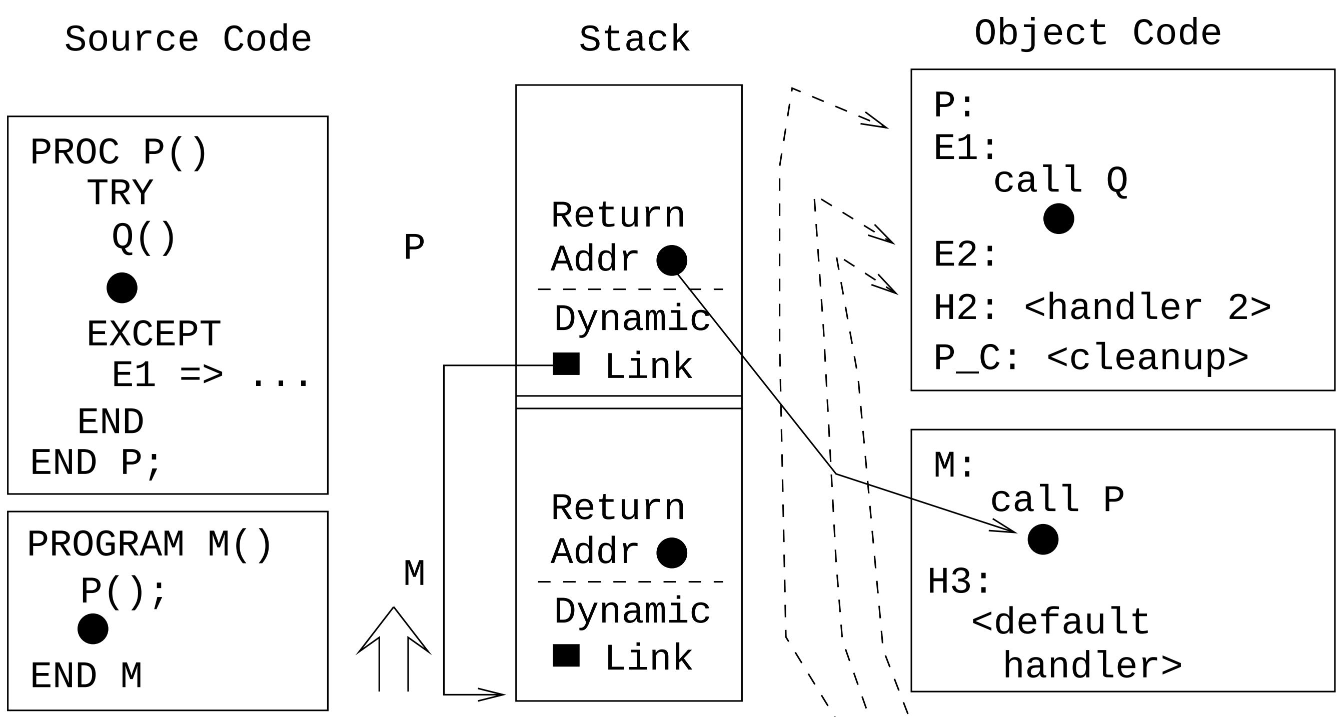
Implementation

- We want 0-overhead exception handling. This means that – unless an exception is raised – there should be no cost associated with the exception handling mechanism.
- We allow raising and handling an exception to be quite slow.
- When an exception is raised we need to be able to
  1. in the current procedure find the exception handler (if any) that encloses the statement that raised the exception, and
  2. rewind the stack (pop activation records) until a procedure with an exception handler is found.

The Range Table

- We build a *RangeTable* at compile-time. It has one entry for each procedure and for each TRY-block.
- Each entry holds four addresses: `pc_high`, `pc_low`, `handler` and `cleanup`.
- [`pc_low`$\cdots$`pc_high`] is the range of addresses for which `handler` is the exception handler.

The Range Table. . .



	(1)	(2)
pc_low:	M	E1
pc_high:	M_end	E2
handler:	H3	H2
cleanup:	/	P_C

Range Table

Unwinding the Stack (Locate)

- Let procedure `S` raise exception `E` at code address `V`. We search the range table to find an entry which covers `V`, i.e. for which `pc_low<=V<=pc_high`.
- Entry `(6)` covers all of procedure `S` (for `S` to `S_end`), and hence `V`. There's no exception handler for this range. We just execute `S`'s cleanup code, `S_C`.
- `S_C` will restore saved registers, etc, and deallocate the activation record.

Unwinding the Stack (Locate)...

Source Code

```
PROC S()
  RAISE E1
END R;
```

Stack

S

```
Return
Addr ●
- - - - - - - - -
Dynamic
■ Link
```

Object Code

```
S:
V: RAISE E1
S_C: <cleanup>
```

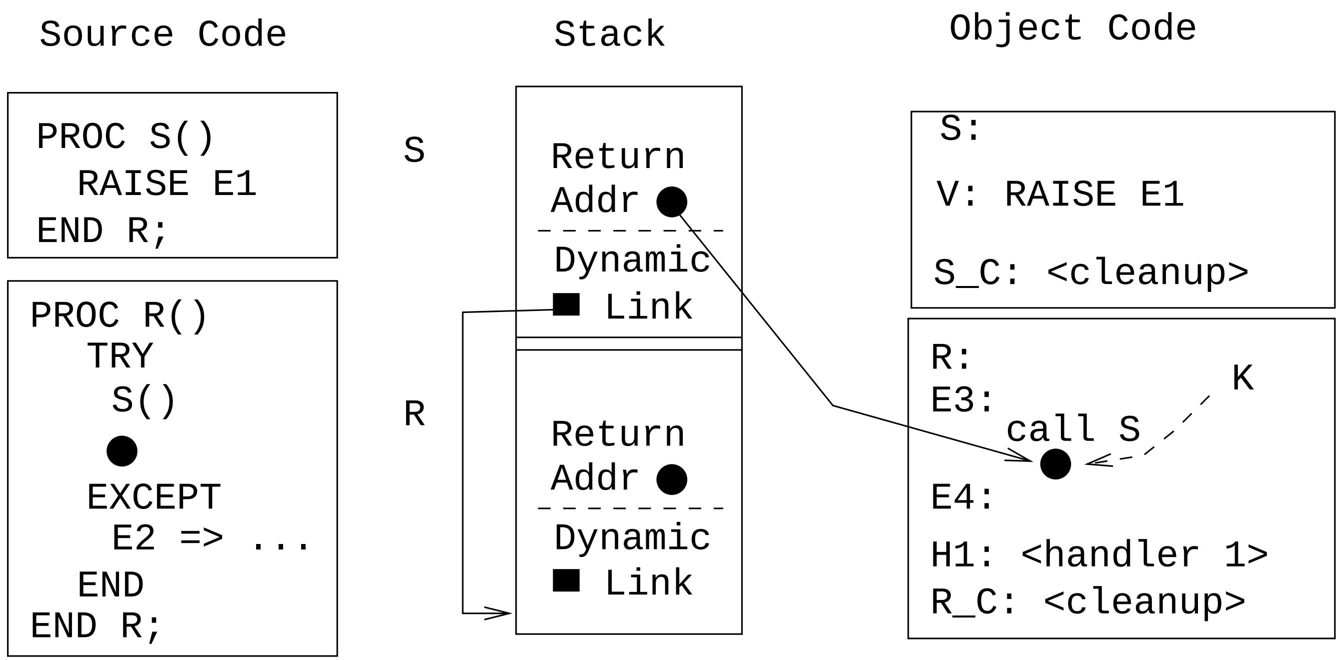
	(5)	(6)
pc_low:	E3	S
pc_high:	E4	S_end
handler:	H1	/
cleanup:	R_C	S_C

Range Table

Unwinding the Stack (Unwind)

- Since `S` didn't have a handler, we must unwind the stack until one is found.
- `S`'s return address is `K`, which is covered by entry `(5)` in the range table. Entry `(5)` has a handler defined (at address `H1`). Run it!

Unwinding the Stack (Unwind)…



	(5)	(6)
pc_low:	E3	S
pc_high:	E4	S_end
handler:	H1	/
cleanup:	R_C	S_C

Range Table

The Exception Handler

- The exception handler itself can be translated as a sequential search.
- If the TRY-EXCEPT-block has no ELSE part, the default action will be to re-raise the exception.

TRY
   S_1;
   RAISE e;
   S_2;
EXCEPT
   E_1 => H_1 |
   E_2 => H_2 |
END;
	$\Longrightarrow$	S_1;
		RAISE e;
		S_2;
		IF $e = E_1$ THEN H_1
		ELSIF $e = E_2$ THEN H_2
		ELSE
		   RAISE e
		ENDIF;

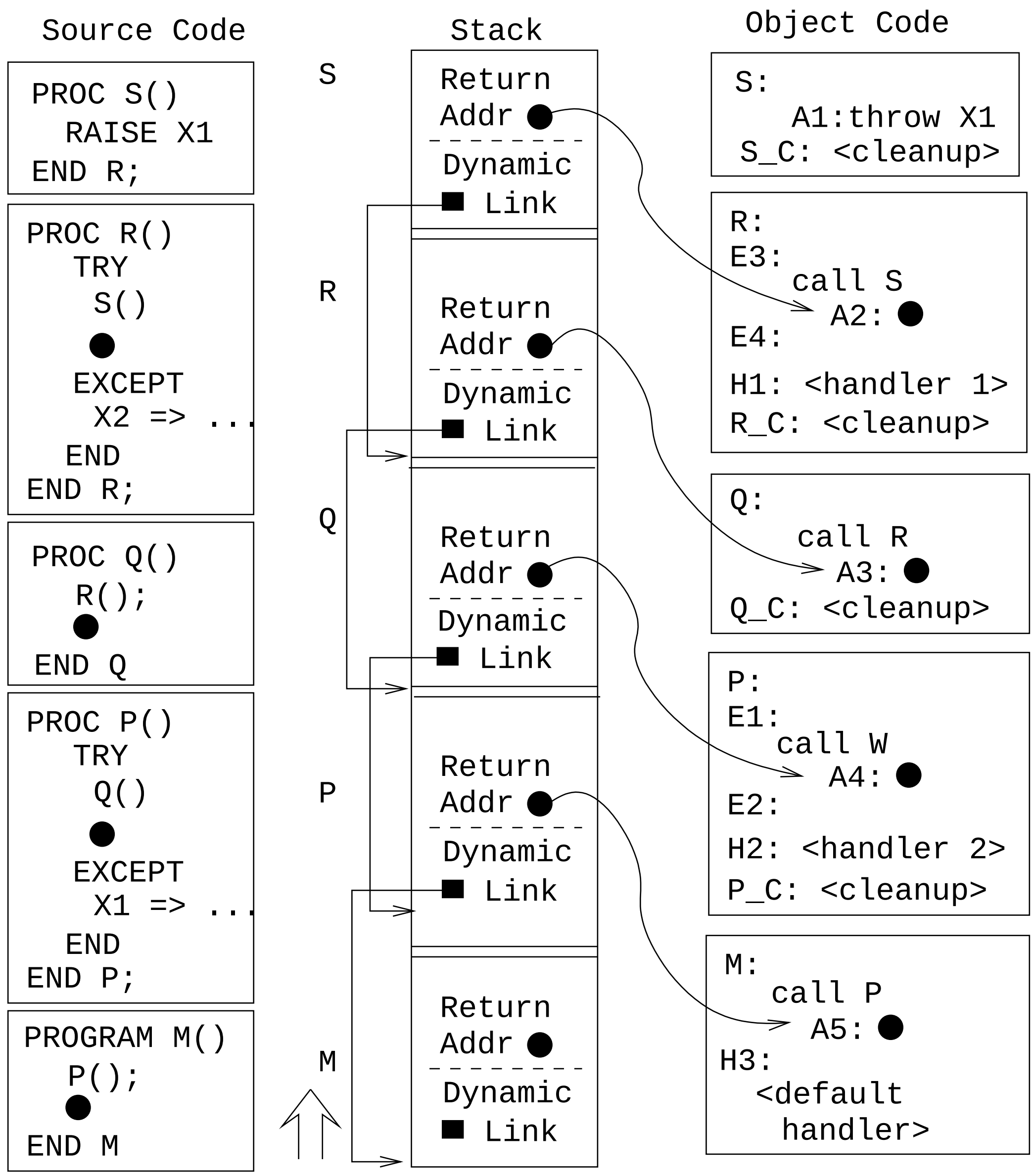
The Algorithm

```
LOOP
   D := The first procedure descriptor (Range Table
         entry) such that D.pc_low <= PC <= D.pc_high;
   IF D.handler = the default handler THEN
      abort and coredump
   ELSIF D.handler ≠ NIL THEN GOTO D.handler;
   ELSE
      Execute the cleanup routine D.cleanup;
      PC := Return address stored in the current frame;
      SP := SP of previous frame;
      FP := FP of previous frame;
   END;
END;
```

Example — Explanation of source code

- Consider the example on the next slide.
- The main program calls procedure `P()`. There is a `<default handler>` defined for the program at address `H3`.
- Procedure `P()` calls `Q()`. Exception `X1` is caught by the handler at address `H2`.
- `Q()` calls `R()`.
- `R()` calls `S()`. Exception `X2` is caught by the handler at address `H1`.
- `S()` throws exception `X1` at address `A1`.

Example



Range Table

	(1)	(2)	(3)	(4)	(5)	(6)
pc_high:	M	E1	P	Q	E3	S
pc_low:	M_end	E2	P_end	Q_end	E4	S_end
handler:	H3	H2	/	/	H1	/
cleanup:	/	P_C	P_C	Q_C	R_C	S_C

Example — Explanation of Actions

- `A1`$\in$`[S,S_end]`, in Range Table entry `(6)`. `(6)` has no handler, so we execute its cleanup routine (`S_C`) and update `PC` to the return address, `A2`.
- Since `A2`$\in$`[E3,E4]` in Range Table entry `(5)`, and `(5).handler==H1`$\neq$`NIL`, we `GOTO H1`. This handler doesn't handle exception `X1`, so it will simply re-raise `X1`.
- `Q()` has no handler, so we execute its cleanup routine (`Q_C`) and propagate the exception to `P()`. I.e. We update `PC` to the return address stored in `Q`'s frame, `A4`.
- Since `A4`$\in$`[E1,E2]` in Range Table entry `(2)`, and `(2).handler=H2`, we `GOTO H2`. This handler catches `X1`. $\Rightarrow$ Done.

Exceptions in C

setjmp/longjmp

- In C, `setjmp/longjmp` can be used to implement exceptional control flow:

```
if (!setjmp(buffer)) {
   /* setjmp returned 0. Protected code.*/
   ...
   longjmp(buffer);
   ...
} else {
   /* setjmp returned 1. Handler code. */
}
```

setjmp/longjmp. . .

- The first time `setjmp` returns 0 and execution continues as normal. When `longjmp` is called it appears as if `setjmp` has returned for the second time, this time returning `1`. The state is now the same as it was when `setjmp` was first called.
- `setjmp`'s buffer argument stores the program's current state, in particular register values.
- Unlike a "real" exception handler, the stack is not rewound nicely. Rather, all stack frames are thrown away. This can lead to problems if not all register values have been saved back in memory. Variables that may be thus affected should be declared `volatile`, i.e. they will always be returned to memory after operated on.

Readings and References

- Read Scott: pp. 441–452
- Drew, Gough, Lederman, *Implementing Zero Overhead Exception Handling*,

  `http://www.dstc.qut.edu.au/~gough/zeroex.ps`.
- Drew, Gough, *Exception handling: Expecting the Unexpected*, Computer Language, Vol 32, No 8, pp. 69–87, 1994.

Summary

- The algorithm we've shown has no overhead (not even one instruction), unless an exception is thrown.
- The major problem that we need to solve is finding the procedure descriptor for a particular stack frame.
- An alternative implementation would be to store a pointer in each frame to the appropriate descriptor. The extra space is negligible, but it would cost 1-2 extra instructions per procedure call.