Introduction

Intermediate Representations

- Some compilers use the AST as the only intermediate representation. Optimizations (code improvements) are performed directly on the AST, and machine code is generated directly from the AST.
- The AST is OK for machine-independent optimizations, such as **inlining** (replacing a procedure call with the called procedure's code).
- The AST is a bit too high-level for machine code generation and machine-dependent optimizations.
- For this reason, some compilers generate a lower level (simpler, closer to machine code) representation from the AST. This representation is used during code generation and code optimization.
Intermediate Code I

Advantages of:
1. Fitting many front-ends to many back-ends,
2. Different development teams for front- and back-end,
3. Debugging is simplified,
4. Portable optimization.

Requirements:
1. Architecture independent,
2. Language independent,
3. Easy to generate,
4. Easy to optimize,
5. Easy to produce machine code from.

A representation which is both architecture and language independent is known as an **UNCOL**, a Universal Compiler Oriented Language.

Intermediate Code II

- UNCOL is the **holy grail** of compiler design – many have search for it, but no-one has found it. Problems:
  1. Programming language semantics differ from one language to another,

- There are several different types of intermediate representations:
  1. Tree-Based.
  2. Graph-Based.
  3. Tuple-Based.
  4. Linear representations.

- All representations contain the same information. Some are easier to generate, some are easy to generate simple machine code from, some are easy to generate **good** code from.

Postfix Notation

**Infix**: \[ b := (a \times 2) + (a \times 2) \]

**Postfix**: \[ b a 2 * a 2 * + := \]

- Postfix notation is a parenthesis free notation for arithmetic expression. It is essentially a linearized representation of an abstract syntax tree.
- In postfix notation an operator appears **after** its operands.
- Very simple to generate, very compact, easy to generate straight-forward machine code from, difficult to generate **good** machine code from.
Tree & DAG Representations

- Trees make good intermediate representations. We can represent the program as a sequence of expression trees. Each assignment, procedure call, or jump becomes one individual tree in the forest.

- **Common Subexpression Elimination** (CSE): Even if the same (sub-) expression appears more than once in a procedure, we should only compute its value once, and save the result for future reference.

- One way of doing this is to build a graph representation, rather than a tree. In the following slides we see how the expression \( a \times 2 \) gets two subtrees in the tree representation and one subtree in the DAG representation.
Tuple Codes

- Another common representation is **three-address code**. It is akin to **assembly code**, but uses an infinite number of **temporaries** (registers) to store the results of operations.
- There are three common realizations of three-address code: **quadruples**, **triples** and **indirect triples**.

Types of 3-Addr Statements:

- \( x := y \ op \ z \): Binary arithmetic or logical operation. Example: Mul, And.
- \( x := \ op \ y \): Unary arithmetic, conversion, or logical operation. Example: Abs, UnaryMinus, Float.
- \( x := y \): Copy statement.
- \( \text{goto } L \): Unconditional jump.

### Three-Address Code II

- \( \text{if } x \ relop y \ \text{goto } L \): Conditional jump. relop is one of \(<, >, \leq, \text{etc.} \) If \( x \ relop y \) evaluates to True, then jump to label L. Otherwise continue with the next tuple.
- \( \text{param } X; \ \text{call } P, n \): Make X the next parameter; make a procedure call to P with \( n \) parameters.
- \( x := y[i] \): Indexed assignment. Set \( x \) to the value in the location \( i \) memory units beyond \( y \).
- \( x := \text{ADDR}(y) \): Address assignment. Set \( x \) to the address of \( y \).
- \( x := \text{IND}(y) \): Indirect assignment. Set \( x \) to the value stored at the address in \( y \).
- \( \text{IND}(x) := y \): Indirect assignment. Set the memory location pointed to by \( x \) to the value held by \( y \).

### Three-Address Code III

- Many three-address statements (particularly those for binary arithmetic) consist of one operator and three addresses (identifiers or temporaries):
  
  \[
  \begin{align*}
  b & := (a \ast 2) + (a \ast 2) \\
  t_1 & := a \ \text{mul} \ 2 \\
  t_2 & := a \ \text{mul} \ 2 \\
  t_3 & := t_1 \ \text{add} \ t_2 \\
  b & := t_3
  \end{align*}
  \]

- There are several ways of implementing three-address statements. They differ in the amount of space they require, how closely tied they are to the symbol table, and how easily they can be manipulated.
- During optimization we may want to move the three-address statements around.
### Three-Address Code IV

**Quadruples:**
- Quadruples can be implemented as an array of records with four fields. One field is the operator.
- The remaining three fields can be pointers to the symbol table nodes for the identifiers. In this case, literals and temporaries must be inserted into the symbol table.

\[
b := (a \times 2) + (a \times 2)
\]

<table>
<thead>
<tr>
<th>Nr</th>
<th>Res</th>
<th>Op</th>
<th>Arg1</th>
<th>Arg2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>t1</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>2</td>
<td>t2</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>3</td>
<td>t3</td>
<td>add</td>
<td>t1</td>
<td>t2</td>
</tr>
<tr>
<td>4</td>
<td>t3</td>
<td>assign</td>
<td>b</td>
<td>t3</td>
</tr>
</tbody>
</table>

### Three-Address Code V

**Triples:**
- Triples are similar to quadruples, but save some space.
- Instead of each three-address statement having an explicit result field, we let the statement itself represent the result.
- We don’t have to insert temporaries into the symbol table.

\[
b := (a \times 2) + (a \times 2)
\]

<table>
<thead>
<tr>
<th>Nr</th>
<th>Op</th>
<th>Arg1</th>
<th>Arg2</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>2</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>3</td>
<td>add</td>
<td>(1)</td>
<td>(2)</td>
</tr>
<tr>
<td>4</td>
<td>assign</td>
<td>b</td>
<td>(3)</td>
</tr>
</tbody>
</table>

### Three-Address Code VI

**Indirect Triples:**
- One problem with triples (“The Trouble With Triples”\(^a\)) is that the around. We may want to do this during optimization.
- We can fix this by adding a level of indirection, an array of pointers

\[
b := (a \times 2) + (a \times 2)
\]

<table>
<thead>
<tr>
<th>Abs</th>
<th>Real</th>
<th>Nr</th>
<th>Op</th>
<th>Arg1</th>
<th>Arg2</th>
</tr>
</thead>
<tbody>
<tr>
<td>(1)</td>
<td>(10)</td>
<td>(11)</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>(2)</td>
<td>(11)</td>
<td>(12)</td>
<td>mul</td>
<td>a</td>
<td>2</td>
</tr>
<tr>
<td>(3)</td>
<td>(12)</td>
<td>(13)</td>
<td>add</td>
<td>(11)</td>
<td>(12)</td>
</tr>
<tr>
<td>(4)</td>
<td>(13)</td>
<td>(14)</td>
<td>:=</td>
<td>b</td>
<td>(13)</td>
</tr>
</tbody>
</table>

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\(^a\)This is a joke. It refers to the famous Star Trek episode “The Trouble With Tribbl
After semantic analysis we traverse the AST and emit the correct intermediate code.

The next slide shows how an expression tree is generated from an AST. The `float` can easily be inserted since all types are known in the AST.

Generating Quadruples I

Each AST node in an expression sub-tree is given an attribute `⇑Place: SymbolT` which represents the name of the identifier or temporary in which the value of the subtree will be computed.

```pascal
PROCEDURE Program (n: Node);
    Decl(n.DeclSeq); Stat(n.StatSeq);
END;

PROCEDURE Decl (n: Node);
    IF n.Kind = ProcDecl THEN
        Decl(n.Locals); Decl(n.Next);
    Stat(n.StatSeq);
    ENDIF
END;

PROCEDURE Stat (n: Node);
    IF n.Kind = Assign THEN
        Expr(n.Des); Expr(n.Expr);
        Emit(n.Des.Place ':=' n.Expr.Place);
    ENDIF
END;

PROCEDURE Expr (n: Node);
    IF n.Kind = Add THEN
        Expr(n.LOP); Expr(n.ROP);
        n.Place := NewTemp();
        Emit(n.Place ':=' n.LOP.Place '+' n.ROP.Place);
    ELSIF n.Kind = VarRef THEN
        n.Place := n.Symbol;
    ENDIF
END;
```

Generating Quadruples II

NewTemp generates a new temporary var.

```pascal
PROCEDURE Stat (n: Node);
    IF n.Kind = Assign THEN
        Expr(n.Des); Expr(n.Expr);
        Emit(n.Des.Place ':=' n.Expr.Place);
    ENDIF
END;

PROCEDURE Expr (n: Node);
    IF n.Kind = Add THEN
        Expr(n.LOP); Expr(n.ROP);
        n.Place := NewTemp();
        Emit(n.Place ':=' n.LOP.Place '+' n.ROP.Place);
    ELSIF n.Kind = VarRef THEN
        n.Place := n.Symbol;
    ENDIF
END;
```
The book uses a convenient way to describe attribute computations, called an **attribute grammar** notation.

We simply combine the abstract syntax notation with the attribute computations that have to be performed at the corresponding AST nodes:

\[
\begin{align*}
A & ::= B \ C \\
& \{ \ C.d := B.c + 1; \\
& \quad A.b := A.a + B.c; \}
\end{align*}
\]

Note that it is not directly obvious from this notation which attributes are synthesized and inherited, and in which order the nodes should be visited. We have to figure this out ourselves!

Now we can rewrite our tuple generator using the new notation:

```plaintext
Assign ::= Des::Expr Expr::Expr
\{ Emit(n.Des.Place ':=' n.Expr.Place); \}
```

```plaintext
Add ::= LOP::Expr ROP::Expr
\{ n.Place := NewTemp(); Emit(n.Place ':=' n.LOP.Place '+' n.ROP.Place); \}
```

```plaintext
Name ::= Ident
\{ n.Place := n.Symbol; \}
```
Building DAGs

From an expression/expression tree such as this one:

\[ a \times (b + c) \]

We might generate this machine code (for some fictitious architecture):

```
LOAD b, r0
LOAD c, r1
ADD r0, r1, r2
LOAD a, r3
MUL r2, r3, r4
```

Can we generate better code from a DAG than a tree?

Example Expression:

\[ [(a + b) \times c + \{(a + b) + e\} \times (e + f)] \times [(a + b) \times c] \]

Tree Representation:

```
+    *
|    |
+    *
|
+    +
|
+    +
```

DAG Representation:

```
+    *
|    |
+    *
|
+    +
|
+    +
```

Generating machine code from the tree yields 21 instructions.

```
LOAD a, r0  ; a
LOAD b, r1  ; b
ADD r0, r1, r2  ; a + b
LOAD c, r0  ; c
MUL r0, r2, r3  ; (a + b) \times c
LOAD a, r0  ; a
LOAD b, r1  ; b
ADD r0, r1, r2  ; a + b
LOAD e, r0  ; e
ADD r2, r0, r4  ; (a + b) + e
LOAD f, r0  ; f
LOAD e, r1  ; e
ADD r0, r1, r0  ; f + e
MUL r4, r0, r4  ; \{(a + b) + e\} \times (e + f)
```
Generating machine code from the DAG yields only 12 instructions.

<table>
<thead>
<tr>
<th>Code from DAG</th>
</tr>
</thead>
<tbody>
<tr>
<td>LOAD a, r0 ; a</td>
</tr>
<tr>
<td>LOAD b, r1 ; b</td>
</tr>
<tr>
<td>ADD r0, r1, r2 ; a + b</td>
</tr>
<tr>
<td>LOAD c, r0 ; c</td>
</tr>
<tr>
<td>MUL r0, r2, r3 ; (a + b) * c</td>
</tr>
<tr>
<td>LOAD e, r4 ; e</td>
</tr>
<tr>
<td>ADD r4, r2, r1 ; (a + b) + e</td>
</tr>
<tr>
<td>LOAD f, r0 ; f</td>
</tr>
<tr>
<td>ADD r0, r4, r0 ; f + e</td>
</tr>
</tbody>
</table>

Repeatedly add subtrees to build DAG. Only add subtrees not already in DAG. Store subtrees in a hash table. This is the \textit{value-number} algorithm.

For every insertion of a subtree, check if \((X \ OP \ Y) \in \text{DAG}\).

\textbf{PROCEDURE} InsertNode (\ OP : Operator; \ L, \ R : Node) : Node;
\begin{align*}
\text{BEGIN} \\
V := \text{hashfunc (OP, L, R)}; \\
N := \text{HashTab.Lookup (V, OP, L, R)}; \\
\text{IF } N = \text{NullNode} \text{ THEN} \\
N := \text{NewNode (OP, L, R)}; \\
\text{HashTab.Insert (V, N)}; \\
\text{END}; \\
\text{RETURN } N; \\
\text{END InsertNode;}
\end{align*}
Read the Tiger book:  
Translation to Intermediate Code  Chapter 7.

Or, read the Dragon book:  
Postfix notation  33  
DAGs & Value Number Alg.  290–293  
Intermediate Languages  463–468, 470–473  
Assignment Statements  478–481

We use an intermediate representation of the program in order to isolate the back-end from the front-end.  
A high-level intermediate form makes the compiler retargetable (easily changed to generate code for another machine). It also makes code-generation difficult.  
A low-level intermediate form make code-generation easy, but our compiler becomes more closely tied to a particular architecture.

Homework  
Translate the program below into quadruples, triples, and a 'sequence of expression trees.'

PROGRAM P;  
VAR X : INTEGER;  
VAR Y : REAL;  
BEGIN  
  X := 1;  
  Y := 5.5;  
  WHILE X < 10 DO  
    Y := Y + FLOAT(X);  
    X := X + 1;  
    IF Y > 10 THEN  
      Y := Y * 2.2;  
    ENDIF;  
  ENDDO;  
END.
Consider the following expression:

\[ ((x \times 4) + y) \times (y + (4 \times x)) + (z \times (4 \times x)) \]

1. Show how the value-number algorithm builds a DAG from the expression (remember that + and \times are commutative).
2. Show the resulting DAG when stored in an array.
3. Translate the expression to postfix form.
4. Translate the expression to indirect triple form.